## FUjiltsu

## FUJITSU Software

## FUJITSU

SSL IIIMPI User's Guide
(Scientific Subroutine Library)

## Preface

This manual describes the functions and usage of the Scientific Subroutine Library II/MPI (SSL II/MPI).

SSL II/MPI provides the computational functionality to efficiently compute large-scale problems on a parallel computer with distributed memory system. The algorithms for parallel processing have been adopted.

This manual consists of two parts.

## Part I General Description

General rules which should be kept in mind when using SSL II/MPI are outlined.

## Part II Usage of Subroutines

The functions and usage of each subroutine are described in alphabetical order of their subroutine names.

Readers of this manual are assumed to be familiar with the MPI system.
For details, refer to the user's guide for MPI.

The asterisks in the table of contents and subroutine list of this manual indicate items added or changed from the previous version.

## Export Controls

Exportation/release of this document may require necessary procedures in accordance with the regulations of your resident country and/or US export control laws.

Date of Publication and Version

| Version | Manual code |
| :--- | :--- |
| March 2021, Version 6.1 | J2UL-2488-02ENZO (01) |
| March 2020, 6th Version | J2UL-2488-02ENZO (00) |
| January 2020, 5th Version | J2UL-2488-01ENZO (00) |
| October 2014, 4th Version | J2UL-1909-01ENZO (00) |
| March 2014, 3rd Version | J2UL-1796-02ENZO (00) |
| February 2012, 2nd Version | - |
| April 2007, 1st Version | - |

## Copyright

Copyright FUJITSU LIMITED 2007-2021

## Update History

| Changes | Location | Version |
| :--- | :--- | :--- |
| The following routine was added. <br> • DS_V3DRCF3 | SSL II/MPI Subroutine List, <br> Usage of Subroutines | 3rd Version |
| • The parameter NW's complemental explanation <br> was added. <br> - Description of Notes 2) was revised. | DS_V3DCFT3 |  |
| The parameter X's explanation was corrected. | DS_V3DRCF |  |
| Rework format | Cover, Preface | 4th Version |
| The following routine was added. <br> D <br> DS_V3DCFT2X | SSL II/MPI Subroutine List, <br> Usage of Subroutines | 5th Version |
| Changed the look according to product upgrades. | - |  |
| Single precision routines are added. | SS_V3DCFT2X, <br> SS_V3DRCF2X | Version 6.1 |
| Correction of a wrong word. | DS_V3DCFT2X, <br> DS_V3DRCF2X |  |

- All rights reserved.
- The information in this manual is subject to change without notice.


## SSL II/MPI Subroutine List

Fourier transforms

| Subroutine name | Item | Page |
| :--- | :--- | :---: |
| DS_V3DCFT | Three-dimensional discrete complex Fourier transforms <br> (mixed radix of 2, 3, 5 and 7, slab decomposition) | 8 |
| DS_V3DCFT2X | Three-dimensional discrete complex Fourier transforms <br> (mixed radix of 2, 3, 5 and 7, pencil decomposition) | 13 |
| DS_V3DCFT3 | Three-dimensional discrete complex Fourier transforms <br> (mixed radix of 2, 3, 5 and 7, volumetric decomposition) | 21 |
| DS_V3DRCF | Three-dimensional discrete real Fourier transforms (mixed <br> radix of 2, 3, 5 and 7, slab decomposition) | 27 |
| DS_V3DRCF2X | Three-dimensional discrete real Fourier transforms (mixed <br> radix of 2, 3, 5 and 7, pencil decomposition) | 33 |
| DS_V3DRCF3 | Three-dimensional discrete real Fourier transforms (mixed <br> radix of 2, 3, 5 and 7, volumetric decomposition) | 41 |
| SS_V3DCFT2X * | Three-dimensional discrete complex Fourier transforms <br> (mixed radix of 2, 3, 5 and 7, pencil decomposition, single <br> precision) | 48 |
| SS_V3DRCF2X * | Three-dimensional discrete real Fourier transforms (mixed <br> radix of 2, 3, 5 and 7, pencil decomposition, single precision) | 56 |

## Contents

## Part I General Description

Chapter 1 Outline ..... 1
Chapter 2 General Rules ..... 3
2.1 Precision of Subroutines ..... 3
2.2 Subroutine Names ..... 3
2.3 Parameters ..... 3
2.4 How to Use SSL II/MPI ..... 3
2.5 Condition Codes ..... 5
Part II Usage of Subroutines
DS_V3DCFT ..... 8
DS_V3DCFT2X ..... 13
DS_V3DCFT3 ..... 21
DS_V3DRCF ..... 27
DS_V3DRCF2X ..... 33
DS_V3DRCF3 ..... 41
SS_V3DCFT2X ..... 48 *
SS_V3DRCF2X ..... 56 *
Appendix
Appendix A References ..... 67
Index ..... IN-1

## Part I <br> General Description

## Chapter 1 Outline

SSL II/MPI is a mathematical software library for parallel execution in a parallel computer with a distributed memory system. It provides the subroutines needed to efficiently compute large-scale problems by parallel processing.

Each SSL II/MPI function is supplied as a subroutine in Fortran. Every subroutine can be referred with a CALL statement.

The functional range, subroutine names, and calling mode of SSL II/MPI are different from those used in the mathematical software library SSL II of the uni-processor version.

## Chapter 2 <br> General Rules

### 2.1 Precision of Subroutines

SSL II/MPI provides subroutines that compute in double-precision, and some routines support the corresponding single-precision.

### 2.2 Subroutine Names

SSL II/MPI consists of user subroutines, that are callable by users and slave subroutines.
Each user subroutine has DS_V or SS_V in the first four characters of its subroutine name.
Each slave subroutine name has DS_U or SS_U in the first four characters.
This manual describes how to use user subroutines.

### 2.3 Parameters

(1) Order of parameters sequence

In general, the order of parameters sequence is the same as that in SSL II:
(input and output parameter list, input parameter list, output parameter list, ICON)
(2) Parameter types

Parameters beginning with I, J, K, L, M, or N are integer type. Parameters beginning with other characters are of the real or complex floating point type. Please refer to the parameter description of each subroutine for the number of bytes of the integer type and the single or double precision of the floating point type.

### 2.4 How to Use SSL II/MPI

(1) The subroutines in SSL II/MPI can be available after the initialization of MPI environment by calling MPI_INIT.

The communicator is specified in the argument of the subroutines. The subroutines make parallel computation in use of the processes belong to the communicator.
The data must be distributed among the processes belong to the communicator.
(2) An example how to use SSL II/MPI
a. Three dimensional complex Fourier transform is considered as an example.

A double precision complex three dimensional array $\mathrm{X}(\mathrm{KX1}, \mathrm{KX} 2, \mathrm{KX} 3 \mathrm{P})$ is allocated in each process. The complex three dimensional data to be transformed is assumed as a matrix of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$.
The (rank+1)-th sub matrix, into which D is partitioned in the third dimension equally by the size KX3P, is stored in the array X on the process of the $\operatorname{rank}(0, \ldots, p-$ 1) ( $p$ is the total number of processes) obtained by MPI_RANK of an MPI subroutine.

```
\(\mathrm{X}(1: \mathrm{N} 1,1: \mathrm{N} 2,1: \mathrm{N} 3 \mathrm{P}) \leftarrow \mathrm{D}(1: \mathrm{N} 1,1: \mathrm{N} 2, \mathrm{~N} 3 \mathrm{~S}: \mathrm{N} 3 \mathrm{E})\) where N3S = KX3P×rank+1,
N3E \(=\operatorname{MIN}(N 3, \mathrm{KX} 3 \mathrm{P} \times(r a n k+1))\), N3P \(=\) MAX(0, N3E-N3S+1).
```

The computation is done in parallel using the distributed data along processes as above.

```
    ** example program **
    use mpi
    implicit real*8 (a-h,o-z)
    parameter (mpn=8)
    parameter (n1=512,n2=n1,n3=n2)
    parameter (kx1=n1+1)
    parameter (kw2p=((n2+mpn-1)/mpn),kx2=kw2p*mpn)
    parameter (kx3p=((n3+mpn-1)/mpn),kw3=kx3p*mpn)
    parameter (nwork=388)
    complex*16 x(kx1,kx2,kx3p),w(kx1,kw2p,kw3),
    $ wc(kx1,kx2,kx3p)
    real*8 dwork(nwork)
    call mpi_init( ierr )
    call mpi_comm_size( mpi_comm_world, nump, ierr )
    call mpi_comm_rank( mpi_comm_world, nop, ierr )
    nop=nop+1
```

    \(i x=1000\)
    ix=ix*nump+nop ! different seed
    do i1=1, kx3p
    call dvrau4(ix, \(x(1,1, i 1)\),
    \$ 2*kx1*kx2,
    \& dwork, nwork, icon)
    enddo
    do i3=1, kx3p
    do i2=1, n2
    do i1=1, n1
    wc(i1,i2,i3)=x(i1,i2,i3)
    enddo
    enddo
    enddo
    isn=1
    call ds_v3dcft(x,kx1,kx2,kx3p,
    \$ n1,n2,n3,w,kw2p,kw3,isn,
\$ mpi_comm_world,icon)
if(icon.ne.0) go to 9000
print*,'icon=',icon
isn=-1
call ds_v3dcft(x, kx1, kx2, kx3p,
\$ n1,n2,n3,w,kw2p,kw3,isn,

```
    $ mpi_comm_world,icon)
        if(icon.ne.0) go to 9000
CC
C
        errorx=0
        iof3=(nop-1)*kx3p
        do i1=1,n1
        do i2=1,n2
        do i3=1,min(kx3p,max(n3-iof3,0))
        errorx=max(dabs(dble(wc(i1,i2,i3))-
    $ dble(x(i1,i2,i3))/n1/n2/n3),errorx)
        errorx=max(dabs(dimag(wc(i1,i2,i3))-
    $ dimag(x(i1,i2,i3))/n1/n2/n3),errorx)
        enddo
        enddo
        enddo
C
    call mpi_allreduce(errorx,errormax,1,mpi_double_precision,
    $ mpi_max,mpi_comm_world,ierr)
C
C
    print*,'-----(',n1,',',n2,',',n3,')-----'
    print*,'error=',errormax
C
    endif
    9000 continue
C
    call mpi_finalize( ierr )
C
    stop
    end
```


## Example 2.1

Example of using an SSL II/MPI routine
Fourier transforms (normal and inverse transforms) for three-dimensional complex data are computed in 8 processes.
b. Parallel computation in threads within each process

Computation in each process can be parallelized in use of threads. The number of threads can be specified in the environment variable OMP_NUM_THREADS.

### 2.5 Condition Codes

The ICON parameter is prepared to indicate the status after the execution of SSL II/MPI.
A value between 0 and 90000 is set as the condition code. The values are classified as shown below depending on whether the result is guaranteed.

Table 2.1 Condition codes

| Code | Meaning | Integrity of the result | Classification |
| :---: | :---: | :---: | :---: |
| 0 | Processing has ended normally. | The results are correct. | Normal |
| 1 to 9999 | Processing has ended normally, but auxiliary information was included. |  |  |
| 10000 to 19999 | Processing has ended with the placing of internal restrictions during execution. | The results are correct on the restrictions. | Warning |
| 20000 to 29999 | Processing was discontinued due to abnormal conditions which had occurred during execution. | The results are not correct. | Abnormal |
| 30000 to 39999 | Processing was discontinued due to invalid input parameter. |  |  |
| 90000 | This routine was deleted from the library. Use the currently supported, superior routine. This comment is also indicated in a message output together with these two routine names to the standard error output file. This message is shown below. | - | - |

Message output at ICON=90000:
JNO0001-S : The SSL II/MPI routine 'AAAAA' no longer available, the
better one 'BBBBB' now recommended for use.
Where,
AAAAA : Name of deleted routine
BBBBB : Name of superior routine

## Part II <br> Usage of Subroutines

## DS_V3DCFT

Three-dimensional discrete complex Fourier transforms. (mixed radices of 2, 3, 5 and 7, slab decomposition)

CALL DS_V3DCFT (X, KX1, KX2, KX3P, N1, N2, N3, W, KW2P, KW3, ISN, COMM, ICON)

## (1) Function

The subroutine DS_V3DCFT performs a three-dimensional complex Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of 2, 3, 5 and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 11_{2 j 3}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3} & =\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
& , k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
& , \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, & \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, & \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j j 3}}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

(2) Parameters

X $\qquad$ Input. Complex data.
Three dimensional complex data is regarded as a matrix data of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$.

The matrix data D is equally partitioned in the third dimension by the size of KX3P into submatrices.

This (rank+1)-th submatrix is stored in an array X on a process of the rank $(0, \ldots, p-1)$ ( $p$ is the total number of processes), which is known in use of the subroutine MPI_COMM_RANK of MPI.
$\mathrm{X}(1: \mathrm{N} 1,1: \mathrm{N} 2,1: \mathrm{N} 3 \mathrm{P}) \leftarrow \mathrm{D}(1: \mathrm{N} 1,1: \mathrm{N} 2$, N3S:N3E), where N3S $=$ KX3P $\times$ rank +1, N3E $=\operatorname{MIN}(N 3, K X 3 P \times(r a n k+1))$, N3P $=$ MAX $(0$, N3E N3S+1).

Output. Transformed complex data.
Resultant transformed three dimensional data of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ are stored into an array X in the same distributed way.

This is a double precision complex three-dimensional array X(KX1, KX2, KX3P).

KX1 $\qquad$ Input. The size of the first dimension of arrays X and $\mathrm{W}(\geq \mathrm{N} 1)$. Integer(INTEGER*4).
KX2 $\qquad$ Input. The size of the second dimension of an array X .
$\mathrm{KX} 2=\mathrm{KW} 2 \mathrm{P} \times p(\geq \mathrm{N} 2)$, and $p$ is the total number of processes. Integer(INTEGER*4).
KX3P .......... Input. The size of the third dimension of an array X.
The size by which the third dimension is equally partitioned ( $\mathrm{KX} 3 \mathrm{P} \times p \geq \mathrm{N} 3$, and $p$ is the total number of processes.).

Integer(INTEGER*4).
N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N2 $\qquad$ Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .

## Integer (INTEGER*4)

W $\qquad$ Work area. This is double precision complex three dimensional array $\mathrm{W}(\mathrm{KX} 1$, KW2P, KW3).
KW2P Input. The size of the second dimension of an array W.
(KW2P $\times p \geq \mathrm{N} 2$, and $p$ is the total number of processes.)
Integer (INTEGER*4)
KW3 $\qquad$ Input. The size of the third dimension of an array W.
$\mathrm{KW} 3=\mathrm{KX} 3 \mathrm{P} \times p$, and $p$ is the total number of processes. Integer(INTEGER*4).

ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
COMM $\qquad$ Input. The communicator indicating a set of processes on which data are distributed and by which computation is done in parallel.

Integer (INTEGER*4).
ICON $\qquad$ Output. Condition code. See Table DS_V3DCFT-1.

Table DS_V3DCFT-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |
| 30000 | N1 $<1, \mathrm{~N} 2<1, \mathrm{~N} 3<1, \mathrm{KX} 1<\mathrm{N} 1, \mathrm{KX} 2 \neq \mathrm{KW} 2 \mathrm{P} \times p$, <br> $\mathrm{KX} 3 \mathrm{P} \times p \neq \mathrm{KW} 3, \mathrm{KX} 2<\mathrm{N} 2, \mathrm{KW} 3<\mathrm{N} 3$, the value <br> of ISN is incorrect. | Processing is discontinued. |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\left.\begin{array}{c}
\begin{array}{c}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1 \\
, k_{3}=0,1, \ldots, n_{3}-1
\end{array} \\
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0 k 3=0}^{n 2-1 n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
\quad, j_{1}=0,1, \ldots, n_{1}-1 \\
\quad, j_{2}=0,1, \ldots, n_{2}-1 \\
\quad, j_{3}=0,1, \ldots, n_{3}-1
\end{array}\right\} \begin{aligned}
& \text { where, } \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right), \\
& \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right) \tag{3.2}
\end{aligned}
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j_{j 2 j}}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) The relation between the size of array $X$ and $W$, and the performance.

The data amounts for communication become minimal when KX1 is close to N 1 , $\mathrm{KX} 2=\mathrm{KW} 2 \mathrm{P} \times p$ is the minimum value satisfying $\mathrm{KX} 2 \geq \mathrm{N} 2$ and $\mathrm{KW} 3=\mathrm{KX} 3 \mathrm{P} \times p$ is the minimum value satisfying $\mathrm{KW} 3 \geq$ N3. $p$ is the total number of processes.
b. Example

Three-dimensional FFT is computed in 8 processes.
do i3=1, kx3p
do i2=1, n2
do i1=1, n1
wc(i1,i2,i3)=x(i1,i2,i3)
enddo
enddo
enddo
C
isn=1
call ds_v3dcft(x, kx1, kx2, kx3p,
\$ n1,n2, n3,w,kw2p,kw3,isn,
\$ mpi_comm_world,icon)
if(icon.ne.0) go to 9000
CC

CC

```
    ** example program **
    use mpi
    implicit real*8 (a-h,o-z)
    parameter (mpn=8)
    parameter (n1=512,n2=n1,n3=n2)
    parameter (kx1=n1+1)
    parameter (kw2p=((n2+mpn-1)/mpn),kx2=kw2p*mpn)
    parameter (kx3p=((n3+mpn-1)/mpn),kw3=kx3p*mpn)
    parameter (nwork=388)
    complex*16 x(kx1,kx2,kx3p),w(kx1,kw2p,kw3),
    $ wc(kx1,kx2,kx3p)
    real*8 dwork(nwork)
    call mpi_init( ierr )
    call mpi_comm_size( mpi_comm_world, nump, ierr )
    call mpi_comm_rank( mpi_comm_world, nop, ierr )
    nop=nop+1
    ix=1000
    ix=ix*nump+nop ! different seed
    do i1=1,kx3p
    call dvrau4(ix,x(1,1,i1),
    $ 2*kx1*kx2,
    & dwork, nwork, icon)
    enddo
    enddo
    * (0)
```

```
    print*,'icon=',icon
CC
C
    isn=-1
    call ds_v3dcft(x,kx1,kx2,kx3p,
    $ n1,n2,n3,w,kw2p,kw3,isn,
    $ mpi_comm_world,icon)
    if(icon.ne.0) go to 9000
cc
c
    errorx=0
    iof3=(nop-1)*kx3p
    do i1=1,n1
    do i2=1,n2
    do i3=1,min(kx3p,max(n3-iof3,0))
    errorx=max(dabs(dble(wc(i1,i2,i3))-
    $ dble(x(i1,i2,i3))/n1/n2/n3),errorx)
    errorx=max(dabs(dimag(wc(i1,i2,i3))
    $ dimag(x(i1,i2,i3))/n1/n2/n3),errorx)
        enddo
        enddo
        enddo
        call mpi_allreduce(errorx,errormax,1,mpi_double_precision,
        $
        mpi_max,mpi_comm_world,ierr)
c
C
    print*,'-----(',n1,',',n2,',',n3,')-----'
    print*,'error=',errormax
C
    endif
    9000 continue
c
    call mpi_finalize( ierr )
c
    stop
    end
```


## DS_V3DCFT2X

Three-dimensional discrete complex Fourier transforms. (mixed radices of 2, 3, 5 and 7, pencil decomposition)

CALL DS_V3DCFT2X (X, KX1, KX2, KX2P, KX3P, Z, KZ1, KZ2, KZ3,
NZ1B, NZ1E, NZ2B, NZ2E, N1, N2, N3, W, NW, ISN, IDIR, COMM2, COMM3, ICON)
(1) Function

The subroutine DS_V3DCFT2X performs a three-dimensional complex Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of 2, 3, 5 and 7.
a. The three-dimensional Fourier transform

When $\left\{x_{j 1 j_{j 23}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3} & =\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
, & k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
, & \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
& , \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, & \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j j 3}}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
\quad, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality using pencil decomposition. The global data of the three-dimensional array is to be distributed among a two-dimensional process grid. The local input array and output array of each process store distinct shapes by dividing the global data in different directions of pencils. This
allows the routine to omit communication for transposing back to the original distribution shape.
(2) Parameters

In the following, the global three dimensional complex data is regarded as an array $\mathrm{D}(\mathrm{N} 1$, N2, N3) virtually, and the process grid is regarded as ND2 $\times$ ND3 shape.

X $\qquad$ Input when IDIR $=1$. Complex data

The local array X stores a subarray of D, which corresponds to a columnwise decomposed part of the global array by dividing the second and third dimensions by ND2 and ND3 respectively.

If each Nm is divisible by $\mathrm{ND} m$ ( $m=2,3$ ), setting a width parameter $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m$ is recommended, then the sizes of the subarray can be (N1,KX2P,KX3P) uniformly.

If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$, setting $\mathrm{KXmP}=\mathrm{Nm} / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than KXmP .

The array X in each process stores the subarray of D as follows:

```
X(1:N1, 1:NX2P, 1:NX3P) \leftarrow D(1:N1, NX2B:NX2E, NX3B:NX3E)
NX2B = KX2P > rank2 + 1
NX2E = MIN(N2, KX2P }\times(rank2 + 1))
NX2P = MAX(0, NX2E - NX2B + 1)
NX3B = KX3P x rank3 + 1
NX3E = MIN(N3, KX3P }\times(\mathrm{ rank3 + 1))
NX3P = MAX(0, NX3E - NX3B + 1),
```

where the rank2 and rank3 are ranks of the process in the communicator COMM2 and COMM3 respectively, which are obtained by subroutine MPI_COMM_RANK of MPI.

The input values are not retained after the calculation.
..........Output when IDIR = -1. Transformed complex data.
Resultant transformed three dimensional data of D(N1, N2, N3) are stored into the array X in the same distributed way as stated above.

This is a double precision complex three-dimensional array X(KX1, KX2, NX3P).

KX1 ............ Input. The size of the first dimension of the array X ( $\geq \mathrm{N} 1$ ). Integer(INTEGER*4).

KX2 $\qquad$ Input. The size of the second dimension of the array $\mathrm{X}(\geq \mathrm{NX} 2 \mathrm{P})$.

KX2P .......... Input. The size by which the second dimension of D is equally partitioned when the data are stored in the array X . (KX2P $\times \mathrm{ND} 2 \geq \mathrm{N} 2$ ) Integer(INTEGER*4)

KX3P .......... Input. The size by which the third dimension of D is equally partitioned when the data are stored in the array X . $(\mathrm{KX} 3 \mathrm{P} \times \mathrm{ND} 3 \geq \mathrm{N} 3)$ Integer(INTEGER*4)

Z $\qquad$ Output when IDIR $=1$. Transformed complex data.

The local array Z stores a subarray of D , which corresponds to a columnwise decomposed part of the global array by dividing the first and second dimensions by ND2 and ND3 respectively.

The array Z in each process stores the subarray of D as follows:
Z(1:NZ1P, 1:NZ2P, 1:N3) $\leftarrow ~ D(N Z 1 B: N Z 1 E, ~ N Z 2 B: N Z 2 E, ~ 1: N 3) ~$
NZ1P $=$ MAX $(0$, NZ1E - NZ1B +1$)$
NZ2P $=$ MAX $(0, N Z 2 E-N Z 2 B+1)$.
......... Input when IDIR = -1. Complex data.
The local array Z stores the subarray of D in the same distributed way as stated above.

This is a double precision complex three-dimensional array Z(KZ1, KZ2, KZ3).
KZ1 ............ Output when $\mathrm{NW}=0$. The recommended size for the first dimension of the array Z.

Input when $N W \neq 0$. The size of the first dimension of the array $Z(\geq$ NZ1E NZ1B+1).

Integer(INTEGER*4).
KZ2 ............ Output when NW $=0$. The recommended size for the second dimension of the array Z .

Input when $N W \neq 0$. The size of the second dimension of the array $Z(\geq$ NZ2E-NZ2B+1).

Integer(INTEGER*4).
KZ3 $\qquad$ Output when NW $=0$. The recommended size of the third dimension of the array Z.

Input when $\mathrm{NW} \neq 0$. The size of the third dimension of the array $\mathrm{Z}(\geq \mathrm{N} 3)$. Integer(INTEGER*4).

NZ1B $\qquad$ Output. The starting index for the first dimension of the global array D. Integer(INTEGER*4).

NZ1E .......... Output. The ending index for the first dimension of the global array D. NZ1B and NZ1E indicate which portion of the first dimension within the global array is stored in the local array Z . Integer(INTEGER*4).

NZ2B $\qquad$ Output. The starting index for the second dimension of the global array D. Integer(INTEGER*4).

NZ2E $\qquad$ Output. The ending index for the second dimension of the global array D. NZ2B and NZ2E indicate which portion of the second dimension within the global array is stored in the local array Z .

Integer(INTEGER*4).
N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)

N2 $\qquad$ Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
W $\qquad$ Work area. This is double precision complex one-dimensional array W(NW).
NW $\qquad$ Input / Output. The size of the work array W.
When NW $=0$ specified, the recommended sizes of NW, KZ1, KZ2, and KZ3 are set respectively, and index information is set to NZ1B, NZ1E, NZ2B, and NZ2E.
Integer (INTEGER*8). (See note 4) in (3), "Comments on use.")
ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
IDIR $\qquad$ Input. The direction of transform between arrays is indicated. IDIR $=1$ for the transform from the array X to the array Z . IDIR $=-1$ for the transform from the array Z to the array X . Integer (INTEGER*4).
COMM2 ...... Input. The MPI communicator that represents a set of processes whose size of the process group is ND2, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 2 ) in (3), "Comments on use.")

Integer (INTEGER*4).
COMM3 ...... Input. The MPI communicator that represents a set of processes whose size of the process group is ND3, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 2 ) in (3), "Comments on use.")

Integer (INTEGER*4).
ICON .......... Output. Condition code.
See Table DS_V3DCFT2X-1.
Integer (INTEGER*4).

Table DS_V3DCFT2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |

Table DS_V3DCFT2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 1000 | NW $=0$ is specified. | The recommended sizes of <br> NW, KZ1, KZ2, and KZ3 <br> are set. <br> Index information is set to <br> NZ1B, NZ1E, NZ2B, and <br> NZ2E. <br> There is no output to array X <br> or Z. |
| 25000 | Too small work area. | Processing is discontinued. |
| 30000 | N1<1, N2<1, N3<1, KX1 < N1, KX2 < NX2P, <br> N2 $>$ KX2P $\times$ ND2, N3 $>$ KX3P $\times$ ND3, or the <br> value of ISN or IDIR is incorrect. |  |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |
| 30100 | COMM2 or COMM3 is incorrect. |  |

## (3) Comments on use

a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1  \tag{3.1}\\
, k_{3}=0,1, \ldots, n_{3}-1 \\
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 33=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j 2 j 3}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) Process shape ND2 and ND3

Note that the performance of this routine may deteriorate when ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.
3) Consistency of parameters among processes

The parameters KX2P,KX3P,N1,N2,N3,NW,ISN, and IDIR needs to have same value respectively among all processes, otherwise the result is not guaranteed.
4) The size of work area $W$

The size of the work array NW needs to be about twice the size of array X or Z to be used as send/receive buffers for MPI communication inside the routine. Note that the parameter NW is an 8-byte integer type.
b. Example

Three-dimensional FFT is computed in $2 \times 3$ processes.

```
** example program **
    use mpi
    implicit real*8 (a-h,o-z)
    parameter (n1=512,n2=n1,n3=n2)
    parameter (nd2=2,nd3=3)
    parameter (kx1=n1)
    parameter (kx2p=((n2+nd2-1)/nd2),kx2=kx2p)
    parameter (kx3p=((n3+nd3-1)/nd3))
    parameter (nwrand=388)
    real*8 dwork(nwrand)
    integer comm2,comm3
    integer*8 nw
    complex*16 x(kx1,kx2,kx3p),wc(kx1,kx2,kx3p)
    complex*16,allocatable :: z(:,:,:),w(:)
    --- prepare sub-communicator ---
    call mpi_init(ierr)
    call mpi_comm_size(mpi_comm_world, nsize, ierr)
    call mpi_comm_rank(mpi_comm_world, nrank, ierr)
    ncolory=nrank/nd2
    call mpi_comm_split(mpi_comm_world,ncolory,nrank,
    & comm2,ierr)
    call mpi_comm_size(comm2, nsize2, ierr)
    call mpi_comm_rank(comm2, nrank2, ierr)
    ncolorz=mod(nrank,nd2)
    call mpi_comm_split(mpi_comm_world,ncolorz,nrank,
    &
                            comm3,ierr)
    call mpi_comm_size(comm3, nsize3, ierr)
    call mpi_comm_rank(comm3, nrank3, ierr)
    if(nsize.ne.nd2*nd3 .or. nsize2.ne.nd2 .or.
    & nsize3.ne.nd3) then
        print*,'nsize=',nsize,nsize2,nsize3
        go to 9000
    endif
    --- prepare test-data ---
    nx2=min(kx2p,max(n2-nrank2*kx2p,0))
    nx3=min(kx3p,max(n3-nrank3*kx3p,0))
    ix=1000
```

```
    ix=ix+nrank ! different seed
    do i3=1,nx3
        do i2=1,n\times2
            call dvrau4(ix,x(1,i2,i3),2*n1,dwork,nwrand,icon)
            do i1=1,n1
                wc(i1,i2,i3)=x(i1,i2,i3)
            enddo
        enddo
        enddo
        --- inquire necessary size ---
        nw=0
        call ds_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
& nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
& comm2,comm3,icon)
    if(icon.ne.1000) then
        print*,'icon=',icon
        go to 9000
    endif
    allocate (z(kz1,kz2,kz3),w(nw))
    print*,'nrank,nrank2,nrank3=',nrank,nrank2,nrank3,
& ' Z-pencil x-range=',nz1b,nz1e,
& ' y-range=',nz2b,nz2e,
& ' z-range=',1,n3
    --- forward FFT ---
    idir=1
    isn=1
    call ds_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
& nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
& comm2,comm3,icon)
    if(icon.ne.0) then
        print*,'icon=',icon
        go to 9000
    endif
    --- backward FFT ---
    idir=-1
    isn=-1
    call ds_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
    & nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
    & comm2,comm3,icon)
    if(icon.ne.0) then
        print*,'icon=',icon
        go to 9000
    endif
    --- check result ---
    errorx=0
    do i3=1,nx3
        do i2=1,n\times2
            do i1=1,n1
            errorx=max(cdabs(wc(i1,i2,i3)-
                        x(i1,i2,i3)/n1/n2/n3),errorx)
            enddo
        enddo
    enddo
```

```
    call mpi_allreduce(errorx,errormax,1,mpi_double_precision,
    & mpi_max,mpi_comm_world,ierr)
    if(nrank.eq.0)then
        print*,'num proc=',nsize
        print*,'nd2,nd3=',nsize2,nsize3
        print*,'-----(',n1,',',n2,',',n3,')-----'
        print*,'error=',errormax
    endif
c
9000 continue
    call mpi_comm_free(comm2,ierr)
    call mpi_comm_free(comm3,ierr)
    call mpi_finalize(ierr)
    stop
    end
```


## DS_V3DCFT3

Three-dimensional discrete complex Fourier transforms. (mixed radices of 2, 3, 5 and 7, volumetric decomposition)

CALL DS_V3DCFT3 (X, KX1, KX2, KX1P, KX2P, KX3P, N1, N2, N3, ND1, ND2, ND3, W, NW, ISN, COMM, ICON)

## (1) Function

The subroutine DS_V3DCFT3 performs a three-dimensional complex Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of 2, 3, 5 and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 1 j_{j 23}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 122 k 3}\right\}$.

$$
\begin{align*}
& n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}=\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
&, k_{1}=0,1, \ldots, n_{1}-1 \\
&, k_{2}=0,1, \ldots, n_{2}-1 \\
&, k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
&, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
&, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
&, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j j 3}}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 33=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality on a massively parallel machine. The global data of the three-dimensional array can be distributed among processes which are regarded as a three-dimensional grid, therefore the volumetric decomposition allows the distribution of work more efficiently than a slabwise decomposition in an environment where massively parallel processes are available.
(2) Parameters

X $\qquad$ Input. Complex data.

When the global three dimensional complex data is regarded as an array $\mathrm{D}(\mathrm{N} 1$, N2, N3) virtually, the local array X stores subarray of D distributed along with the shape of the process grid.

If each $\mathrm{N} m$ is divisible by $\mathrm{NDm}(m=1,2,3)$, setting a width parameter $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m$ is recommended, then the sizes of the subarray can be (KX1P,KX2P,KX3P) uniformly.

If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$, setting $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than KXmP .

The array X in each process stores the subarray of D as follows: X (1:N1P, 1:N2P, 1:N3P) $\leftarrow$ D(N1S:N1E, N2S:N2E, N3S:N3E),
N1S $=$ KX1P $\times$ rank1 +1 ,
$\mathrm{N} 1 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 1, \mathrm{KX1P} \times($ rank1 +1$))$,
N1P = MAX(0, N1E - N1S + 1),
N2S $=$ KX2P $\times$ rank2 +1 ,
$\mathrm{N} 2 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 2, \mathrm{KX} 2 \mathrm{P} \times($ rank2 +1$)$ ),
$\mathrm{N} 2 \mathrm{P}=\mathrm{MAX}(0, \mathrm{~N} 2 \mathrm{E}-\mathrm{N} 2 \mathrm{~S}+1)$,
N3S $=$ KX3P $\times$ rank3 +1 ,
N3E $=\operatorname{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times($ rank3 +1$))$, N3P $=$ MAX $(0$, N3E - N3S +1$)$.

The rank1,rank2 and rank3 are coordinates of the process grid calculated as follows from the rank value, which is obtained by subroutine MPI_COMM_RANK of MPI.
rank1 $=\bmod ($ rank, ND1),
rank2 $=\bmod ($ rank/ND1, ND2),
rank3 $=\operatorname{rank} /(\mathrm{ND} 1 \times \mathrm{ND} 2)$.
Output. Transformed complex data.
Resultant transformed three dimensional data of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ are stored into the array X in the same distributed way.

This is a double precision complex three-dimensional array X(KX1, KX2, KX3P).

KX1 ............ Input. The size of the first dimension of the array X ( $\geq \mathrm{KX} 1 \mathrm{P}$ ).
Integer(INTEGER*4).
KX2 $\qquad$ Input. The size of the second dimension of the array $\mathrm{X}(\geq \mathrm{KX} 2 \mathrm{P}$ ). Integer(INTEGER*4).

KX1P ........... Input. The size by which the first dimension is equally partitioned.(KX1P $\times$ ND1 $\geq$ N1) Integer(INTEGER*4).

KX2P $\qquad$ Input. The size by which the second dimension is equally partitioned.(KX2P $\times$ ND2 $\geq$ N2) Integer(INTEGER*4).

KX3P $\qquad$ Input. The size by which the third dimension is equally partitioned (KX3P $\times$ ND $3 \geq$ N3)

## Integer(INTEGER*4).

N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N 2 ............... Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3
Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
ND1 $\qquad$ Input. The number of processes by which the first dimension is partitioned.
Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
ND2 $\qquad$ Input. The number of processes by which the second dimension is partitioned.

Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
ND3 $\qquad$ Input. The number of processes by which the third dimension is partitioned.
Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
W $\qquad$ Work area. This is double precision complex one-dimensional array W(NW).

NW $\qquad$ Input. The size of the work array $\mathrm{W}(\mathrm{NW} \geq \operatorname{MAX}(\mathrm{KX} 1 \times \mathrm{ND} 1, \mathrm{KX} 2 \mathrm{P} \times \mathrm{ND} 2$, $\mathrm{KX} 3 \mathrm{P} \times$ ND3 $) \times 3$ ). It is recommended to specify a sufficiently large size for efficiency. (See note 2) in (2), "Comments on use.")

## Integer (INTEGER*4)

ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
COMM ....... Input. The communicator indicating a set of processes on which data are distributed and by which computation is done in parallel.
Integer (INTEGER*4).
ICON .......... Output. Condition code.
See Table DS_V3DCFT3-1.
Integer (INTEGER*4).

Table DS_V3DCFT3-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |
| 25000 | Too small work area. | Processing is discontinued. |
| 30000 | N1<1, N2<1, N3<1, KX1 $<$ KX1P, KX2 $<$ <br> KX2P, N1 $>$ KX1P $\times$ ND1, N2 $>$ KX2P $\times$ <br> ND2, N3 $>$ KX3P $\times$ ND3, or the value of ISN <br> is incorrect. |  |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |
| 30100 | ND1 $\times$ ND2 $\times$ ND3 is not equal to total <br> processes. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1  \tag{3.1}\\
, k_{3}=0,1, \ldots, n_{3}-1 \\
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j_{j 2 j}}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) The size of work area W

The size of the work array determines partition sizes for transferring data among nodes and calculation on each node in this routine. Setting the size NW of work area W much larger than MAX $(\mathrm{KX} 1 \times \mathrm{ND} 1, \mathrm{KX} 2 \mathrm{P} \times \mathrm{ND} 2, \mathrm{KX} 3 \mathrm{P} \times \mathrm{ND} 3) \times 3 \times$ (number of threads in a process) is recommended. For example, setting NW > 500,000 is expected to be efficient when the assigned cache size to the process is 8 MB and the array X is partitionable by that size of the work area.
3) Parameters ND1, ND2 and ND3

When using volumetric decomposition, it is recommended to adjust ND1,ND2 and ND3 to be about the comparable value of cube root of the number of total process ND $1 \times$ ND $2 \times$ ND3 for overall efficiency of transferring data. Note that the performance of this routine may deteriorate when ND1,ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.
Additionally, this routine can be used for slabwise decomposition or 2dimensional decomposition also by setting any of ND1,ND2 or ND3 to 1, when user's program exploits specific decomposition or available shape of the process grid on a system is limited.
4) Consistency of parameters among processes

The parameters KX1,KX2,KX1P,KX2P,KX3P,N1,N2,N3,ND1,ND2,ND3,NW and ISN needs to have same value respectively among all processes, otherwise the result is not guaranteed.
b. Example

Three-dimensional FFT is computed in $2 \times 2 \times 2$ processes.

C
call mpi_init( ierr )
call mpi_comm_size( mpi_comm_world, nump, ierr )
call mpi_comm_rank( mpi_comm_world, nop, ierr )
nrank1=mod (nop, nd1)
nrank2=mod(nop/nd1, nd2)
nrank3=nop/(nd1*nd2)
$i x=1000$
ix=ix*nump+nop ! different seed
do i1=1, kx3p
call dvrau4(ix, $x(1,1, i 1)$,
\$ 2*k×1*kx2,
\& dwork, nwork, icon)
enddo
C
do i3=1, kx3p
do i2=1, kx2p

```
    do i1=1, kx1p
    wc(i1,i2,i3)=x(i1,i2,i3)
    enddo
    enddo
    enddo
    isn=1
    call ds_v3dcft3(x,kx1,kx2,kx1p,kx2p,kx3p,
    $ n1,n2,n3,nd1,nd2,nd3,w,nw,isn,
    $ mpi_comm_world,icon)
    if(icon.ne.0) go to 9000
    isn=-1
    call ds_v3dcft3(x,kx1,kx2,kx1p,kx2p,kx3p,
    $ n1,n2,n3,nd1,nd2,nd3,w,nw,isn,
    $ mpi_comm_world,icon)
    if(icon.ne.0) go to 9000
cc
c
c
call mpi_allreduce(errorx,errormax,1,mpi_double_precision,
$
    mpi_max,mpi_comm_world,ierr)
if(nop.eq.1)then
    print*,'----(',n1,',',n2,',',n3,')-----'
    print*,'error=',errormax
    endif
    9000 continue
C
    call mpi_finalize( ierr )
c
    stop
end
```


## DS_V3DRCF

Three-dimensional discrete real Fourier transforms. (mixed radices of 2, 3, 5 and 7, slab decomposition)

CALL DS_V3DRCF(X, KX1, KX2, KX3P, N1, N2, N3, W, KW2P, KW3, ISIN, ISN, COMM, ICON)

## (1) Function

The subroutine DS_V3DRCF performs a three-dimensional real Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of the three-dimensional array $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of $2,3,5$ and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 1 j_{j 23}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 112 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}= & \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1 r} \omega_{n 2}^{-j 2 k 2 r} \omega_{n 3}^{-j 3 k 3 r} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
& , k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
& , \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
& , \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
& , \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right) \\
& r=1 \text { or } r=-1
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j 2 j 3}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1 r} \omega_{n 2}^{j 2 k 2 r} \omega_{n 3}^{j 3 k 3 r} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right) \\
r=1 \text { or } r=-1
\end{gather*}
$$

## (2) Parameters

X $\qquad$ Input/Output. Three-dimensional real data.
Three dimensional real data is regarded as a matrix data of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$.

The matrix data D is equally partitioned in the third dimension by the size of KX3P into submatrices.

This (rank+1)-th submatrix is stored in an array X on a process of the rank $(0, \ldots, p-1)$ ( $p$ is the total number of processes), which is known in use of the subroutine MPI_COMM_RANK of MPI.
$\mathrm{X}(1: \mathrm{N} 1,1: \mathrm{N} 2,1: \mathrm{N} 3 \mathrm{P}) \leftarrow \mathrm{D}(1: \mathrm{N} 1,1: \mathrm{N} 2, \mathrm{~N} 3 \mathrm{~S}: \mathrm{N} 3 \mathrm{E})$
Where N3S $=\mathrm{KX} 3 \mathrm{P} \times r a n k+1, \mathrm{~N} 3 \mathrm{E}=\operatorname{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times(r a n k+1))$, N3P $=$ MAX (0, N3E-N3S+1).

For the real to complex transform (ISN = 1), data is input; for the complex to real transform (ISN =-1), data is output.

Output/input. The real and imaginary parts of the transformed complex data.
For the real to complex transform (ISN = 1), data is output; for the complex to real transform (ISN =-1), data is input.

The complex data $\mathrm{CD}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ obtained from real data $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ by Fourier transform has the complex conjugate relation so the about a half of the first dimension $(1 \sim \mathrm{~N} 1 / 2+1)$ is used to store the complex data.
(See note 2) in (3), "Comments on use.")
Regarding an array X as $\mathrm{X}(2, \mathrm{KX} 1 / 2, \mathrm{KX} 2, \mathrm{KX} 3 \mathrm{P})$, the real and imaginary parts are stored in $\mathrm{X}(1,1: \mathrm{N} 1 / 2+1,1: \mathrm{N} 2,1: \mathrm{N} 3)$ and $\mathrm{X}(2,1: \mathrm{N} 1 / 2+1,1: \mathrm{N} 2,1: \mathrm{N} 3)$ respectively.

This is a double precision real three dimensional array X(KX1, KX2, KX3P).
KX1 $\qquad$ Input. The size of the first dimension of array X and $\mathrm{W}(\geq 2 \times \mathrm{N} 1 / 2+1)$. KX1 must be even.

Integer(INTEGER*4).
KX2 $\qquad$ Input. The size of the second dimension of an array X.
$\mathrm{KX} 2=\mathrm{KW} 2 \mathrm{P} \times p$, where $p$ is the total number of processes $(\geq \mathrm{N} 2)$. Integer(INTEGER*4).

KX3P $\qquad$ Input. The size of the third dimension of an array X .
The size by which the third dimension is equally partitioned (KX3P $\times p \geq \mathrm{N} 3$, and $p$ is the total number of processes.).

Integer(INTEGER*4).
N1 $\qquad$ Input. The length $n_{1}$ of real data in the first dimension to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N2 $\qquad$ Input. The length $n_{2}$ of real data in the second dimension to be transformed. $n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 . Integer (INTEGER*4)

N3 $\qquad$ Input. The length $n_{3}$ of real data in the third dimension to be transformed. $n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 . Integer (INTEGER*4)

W $\qquad$ Work area. This is double precision real three dimensional array W(KX1, KW2P, KW3).

KW2P ......... Input. The size of the second dimension of an array W.
(KW2P $\times p \geq \mathrm{N} 2$, and $p$ is the total number of processes.)
Integer (INTEGER*4)
KW3 $\qquad$ Input. The size of the third dimension of an array W.
$\mathrm{KW} 3=\mathrm{KX3P} \times p$, and $p$ is the total number of processes.
Integer(INTEGER*4).
ISIN $\qquad$ Input. The direction of the transformation.
$r=1$ for 1.
$r=-1$ for -1 .
Integer (INTEGER*4)
ISN .......... ... Input. Either the transform or the inverse transform is indicated.
ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
COMM ....... Input. The communicator indicating a set of processes on which data are distributed and by which computation is done in parallel.

Integer (INTEGER*4).
ICON .......... Output. Condition code.
See Table DS_V3DRCF-1.

Table DS_V3DRCF-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |
| 30000 | $\mathrm{KX}<2 \times(\mathrm{N} 1 / 2+1), \mathrm{KX} 1$ is not even, <br> $\mathrm{KX} 2 \neq \mathrm{KW} 2 \mathrm{P} \times p, \mathrm{KX} 3 \mathrm{P} \times p \neq \mathrm{KW} 3, \mathrm{KX} 2<\mathrm{N} 2$, <br> $\mathrm{KW} 3<\mathrm{N} 3, \mathrm{~N} 1<1, \mathrm{~N} 2<1, \mathrm{~N} 3<1, \mathrm{ISIN} \neq 1,-1$, <br> $\mathrm{ISN} \neq 1,-1$. | Processing is discontinued. |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |

(3) Comments on use
a Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1  \tag{3.1}\\
, k_{2}=0,1, \ldots, n_{2}-1 \\
, k_{3}=0,1, \ldots, n_{3}-1 \\
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 112 k 3}\right\}$ or $\left\{x_{j 1 j 2 j 3}\right\}$ corresponding to the left term of (3.1) or (3.2), respectively. The normalization of the results may be required.
2) The results of the three-dimensional real Fourier transform has the following complex conjugate relation (indicated by ${ }^{-}$).

$$
\begin{equation*}
\alpha_{k 1 k 2 k 3}=\overline{\alpha_{n 1-k 1 n 2-k 2 n 3-k 3}} \tag{3.3}
\end{equation*}
$$

The remainder of the data is obtained from data in $k_{1}=0, \ldots, n_{1}-1, k_{2}=0, \ldots, n_{2-}$ 1 , and $k_{3}=0, \ldots, n_{3} / 2$.

Example
Three-dimensional real FFT is computed in 8 processes.
c
call mpi_init( ierr )
call mpi_comm_size( mpi_comm_world, nump, ierr )
call mpi_comm_rank( mpi_comm_world, nop, ierr )
nop=nop+1
$i x=1000$

```
    iof3=(nop-1)*kx3p
```

    error=0
    do i1=1,n1
    do i2=1,n2
    do i3=1, min \((k x 3 p, \max (n 3-i o f 3,0))\)
    error=max(dabs(wc(i1,i2,i3)-
    \$ \(\quad x(i 1, i 2, i 3) / n 1 / n 2 / n 3)\), error \()\)
    enddo
    enddo
    enddo
    call mpi_allreduce(error,errormax,1,mpi_double_precision,
    \$
                                    mpi_max,mpi_comm_world,ierr)
    C
C

print*,'error=',errormax
c
endif
9000 continue
C
call mpi_finalize( ierr )

C
stop
end

## DS_V3DRCF2X

Three-dimensional discrete real Fourier transforms. (mixed radices of 2, 3, 5 and 7, pencil decomposition)

CALL DS_V3DRCF2X (X, KX1, KX2, KX2P, KX3P, Z, KZ1, KZ2, KZ3,
NZ1B, NZ1E, NZ2B, NZ2E, N1, N2, N3, W, NW, ISN, IDIR, COMM2, COMM3, ICON)

## (1) Function

The subroutine DS_V3DRCF2X performs a three-dimensional real Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of 2, 3, 5 and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 1 j_{j 23}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3} & =\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
, & k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
, & \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
& , \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, & \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j} j 3}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality using pencil decomposition. The global data of the three-dimensional array is to be distributed among a two-dimensional process grid. The local input array and output array of each process store distinct shapes by dividing the global data in different directions of pencils. This
allows the routine to omit communication for transposing back to the original distribution shape.
(2) Parameters

In the following, the global three dimensional real data is regarded as an array $\mathrm{DR}(\mathrm{N} 1, \mathrm{~N} 2$, N 3 ), and the transformed global three dimensional complex data is regarded as an array $\mathrm{DC}(\mathrm{N} 1 / 2+1, \mathrm{~N} 2, \mathrm{~N} 3)$ virtually, and the process grid is regarded as ND2 $\times$ ND3 shape.

X $\qquad$ Input when IDIR $=1$. Real data.
The local array X stores a subarray of DR , which corresponds to a columnwise decomposed part of the global array by dividing the second and third dimensions by ND2 and ND3 respectively.

If each $\mathrm{N} m$ is divisible by $\mathrm{NDm}(m=2,3)$, setting width parameters $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m$ is recommended, then the data size settled in the subarray can be (N1,KX2P,KX3P) uniformly.

If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$, setting $\mathrm{KX} m \mathrm{P}=\mathrm{Nm} / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than KXmP .

The array X in each process stores the subarray of DR as follows:
X(1:N1, 1:NX2P, 1:NX3P) $\leftarrow$ DR(1:N1, NX2B:NX2E, NX3B:NX3E),
NX2B $=$ KX2P $\times$ rank2 +1
$\mathrm{NX} 2 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 2, \mathrm{KX} 2 \mathrm{P} \times($ rank2 +1$))$
NX2P $=$ MAX $(0$, NX2E $-N X 2 B+1)$
NX3B $=$ KX3P $\times$ rank3 +1
NX3E $=\operatorname{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times($ rank3 +1$))$
NX3P = MAX (0, NX3E - NX3B + 1),
where, rank2 and rank3 are the ranks of the process in the communicator COMM2 and COMM3 respectively, which are obtained by subroutine MPI_COMM_RANK of MPI.

The input values are not retained after the calculation.
..........Output when IDIR = -1. Transformed real data.
Resultant transformed three dimensional real data from DC(N1/2+1, N2, N3) are stored into the array X in the same distributed way as stated above.

This is a double precision real three-dimensional array X(KX1, KX2, KX3P).
KX1 ............Input. The size of the first dimension of the array X. KX1 must be even. ( $\geq \mathrm{N} 1$ ). Integer(INTEGER*4).

KX2 $\qquad$ Input. The size of the second dimension of the array $\mathrm{X}(\geq \mathrm{NX} 2 \mathrm{P})$. Integer(INTEGER*4).
KX2P $\qquad$ Input. The size by which the second dimension of DR is equally partitioned. (KX2P $\times$ ND $2 \geq$ N2) Integer(INTEGER*4).
KX3P .......... Input. The size by which the third dimension of DR is equally partitioned. (KX3P $\times$ ND $3 \geq$ N3) Integer(INTEGER*4)

Z $\qquad$ Output when IDIR $=1$. Transformed complex data.

The complex data obtained from real data DR(N1, N2, N3) by Fourier transform has the complex conjugate relation so the about a half of the first dimension is used to store. The local array Z stores a subarray of $\mathrm{DC}(\mathrm{N} 1 / 2+1$, N2, N3), which corresponds to a columnwise decomposed part of the global array by dividing the first and second dimensions by ND2 and ND3 respectively.

The array Z in each process stores the subarray of DC as follows: Z(1:NZ1P, 1:NZ2P, 1:N3) $\leftarrow$ DC(NZ1B:NZ1E, NZ2B:NZ2E, 1:N3)
NZ1P $=$ MAX $(0$, NZ1E - NZ1B +1$)$ NZ2P $=$ MAX $(0, N Z 2 E-N Z 2 B+1)$.
......... Input when IDIR $=-1$. Complex data.
The local array Z stores the subarray of DC in the same distributed way as stated above.

This is a double precision complex three-dimensional array Z(KZ1, KZ2, KZ3).
(See note 2) in (3), "Comments on use.")
KZ1 $\qquad$ Output when NW $=0$. The recommended size for the first dimension of the array Z.

Input when NW $\neq 0$. The size of the first dimension of the array $\mathrm{Z}(\geq$ NZ1E NZ1B+1).

Integer(INTEGER*4).
KZ2 ............ Output when NW $=0$. The recommended size for the second dimension of the array Z.

Input when $N W \neq 0$. The size of the second dimension of the array $\mathrm{Z}(\geq$ NZ2E-NZ2B+1).

Integer(INTEGER*4).
KZ3 $\qquad$ Output when NW $=0$. The recommended size of the third dimension of the array Z.

Input when $\mathrm{NW} \neq 0$. The size of the third dimension of the array $\mathrm{Z}(\geq \mathrm{N} 3)$. Integer(INTEGER*4).
NZ1B $\qquad$ Output. The starting index for the first dimension of the global array DC. Integer(INTEGER*4).

NZ1E .......... Output. The ending index for the first dimension of the global array DC. NZ1B and NZ1E indicate which portion of the first dimension within the global array is stored in the local array Z .

Integer(INTEGER*4).
NZ2B $\qquad$ Output. The starting index for the second dimension of the global array DC. Integer(INTEGER*4).

NZ2E .......... Output. The ending index for the second dimension of the global array DC. NZ2B and NZ2E indicate which portion of the second dimension within the global array is stored in the local array Z .

Integer(INTEGER*4).
N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 . Integer (INTEGER*4)

N2 ............... Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed. $n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 . Integer (INTEGER*4)

N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
W $\qquad$ Work area. This is double precision complex one-dimensional array W(NW).

NW $\qquad$ Input / Output. The size of the work array W.

When NW = 0 specified, the recommended sizes of NW, KZ1, KZ2, and KZ3 are set respectively, and the index information is set to NZ1B, NZ1E, NZ2B, and NZ2E.

Integer (INTEGER*8). (See note 5) in (3), "Comments on use.")
ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
IDIR $\qquad$ Input. The direction of transform between arrays is indicated.

IDIR = 1: from the real array X to the complex array Z .
IDIR $=-1$ : from the complex array Z to the real array X .
Integer (INTEGER*4).
COMM2 ...... Input. The MPI communicator that represents a set of processes whose size of the process group is ND2, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 3 ) in (3), "Comments on use.") Integer (INTEGER*4).

COMM3 ...... Input. The MPI communicator that represents a set of processes whose size of the process group is ND3, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 3 ) in (3), "Comments on use.") Integer (INTEGER*4).

ICON $\qquad$ Output. Condition code.

See Table DS_V3DRCF2X-1. Integer (INTEGER*4).

Table DS_V3DRCF2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |

Table DS_V3DRCF2X-1 Condition codes

| Code | Meaning | Processing |
| :---: | :---: | :---: |
| 0 | No error |  |
| 1000 | $\mathrm{NW}=0$ is specified | The recommended sizes of NW, KZ1, KZ2 and KZ3 are set. <br> Index information is set to NZ1B, NZ1E, NZ2B, and NZ2E. <br> There is no output to array X or Z . |
| 25000 | Too small work area. | Processing is discontinued. |
| 30000 | $\begin{aligned} & \mathrm{N} 1<1, \mathrm{~N} 2<1, \mathrm{~N} 3<1, \\ & \mathrm{KX} 1<\mathrm{N} 1, \mathrm{KX} 1 \text { is not even, } \\ & \mathrm{KX} 2<\mathrm{NX} 2 \mathrm{P}, \\ & \mathrm{~N} 2>\mathrm{KX} 2 \mathrm{P} \times \mathrm{ND} 2, \mathrm{~N} 3>\mathrm{KX} 3 \mathrm{P} \times \mathrm{ND} 3, \\ & (\mathrm{~N} 1 / 2+1) \times 2>\mathrm{KZ} 1 \times \mathrm{ND} 2, \\ & \mathrm{I} N \mathrm{~N} \neq 1,-1, \text { IDIR } \neq 1,-1 . \end{aligned}$ |  |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |
| 30100 | COMM2 or COMM3 is incorrect. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1  \tag{3.1}\\
, k_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j 2 j 3}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) The results of the three-dimensional real Fourier transform has the following complex conjugate relation (indicated by ${ }^{-}$).

$$
\begin{equation*}
\alpha_{k 1 k 2 k 3}=\overline{\alpha_{n 1-k 1 n 2-k 2 n 3-k 3}} \tag{3.3}
\end{equation*}
$$

The remainder of the data is obtained from data in $k_{1}=0, \ldots, n_{1}-1, k_{2}=0, \ldots, n_{2-}$ 1 , and $k_{3}=0, \ldots, n_{3} / 2$.
3) Process shape ND2 and ND3

Note that the performance of this routine may deteriorate when ND1,ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.
4) Consistency of parameters among processes

The parameters KX2P, KX3P, N1, N2, N3, NW, ISN and IDIR needs to have same value respectively among all processes, otherwise the result is not guaranteed.
5) The size of work area W

The size of the work array NW needs to be about twice the size of array X or Z to be used as send/receive buffers for MPI communication inside the routine. Note that the parameter NW is an 8-byte integer type.
b. Example

Three-dimensional FFT is computed in $2 \times 3$ processes.

```
    ** example program **
    use mpi
    implicit real*8 (a-h,o-z)
    parameter (n1=512,n2=n1,n3=n2)
    parameter (nd2=2,nd3=3)
    parameter (n1c=n1/2+1, kx1=n1c*2)
    parameter (kx2p=((n2+nd2-1)/nd2),kx2=kx2p)
    parameter (kx3p=((n3+nd3-1)/nd3))
    parameter (nwrand=388)
    real*8 dwork(nwrand)
    integer comm2,comm3
    integer*8 nw
    real*8 x(kx1,kx2,kx3p),wc(kx1,kx2,kx3p)
    complex*16,allocatable :: z(:,:,:)
    real*8,allocatable :: w(:)
    --- prepare sub-communicator ---
    call mpi_init(ierr)
    call mpi_comm_size(mpi_comm_world, nsize, ierr)
    call mpi_comm_rank(mpi_comm_world, nrank, ierr)
    ncolory=nrank/nd2
    call mpi_comm_split(mpi_comm_world,ncolory,nrank,
& comm2,ierr)
    call mpi_comm_size(comm2, nsize2, ierr)
    call mpi_comm_rank(comm2, nrank2, ierr)
```

```
    ncolorz=mod(nrank,nd2)
    call mpi_comm_split(mpi_comm_world,ncolorz,nrank,
    & comm3,ierr)
    call mpi_comm_size(comm3, nsize3, ierr)
    call mpi_comm_rank(comm3, nrank3, ierr)
    if(nsize.ne.nd2*nd3 .or. nsize2.ne.nd2 .or.
& nsize3.ne.nd3) then
        print*,'nsize=',nsize,nsize2,nsize3
        go to 9000
    endif
    --- prepare test-data ---
    nx2=min(kx2p,max(n2-nrank2*kx2p,0))
    nx3=min(kx3p,max(n3-nrank3*kx3p,0))
    ix=1000
    ix=ix+nrank ! different seed
    do i3=1,nx3
        do i2=1,n\times2
            call dvrau4(ix,x(1,i2,i3),n1,dwork,nwrand,icon)
            do i1=1,n1
                    wc(i1,i2,i3)=x(i1,i2,i3)
            enddo
        enddo
    enddo
    --- inquire necessary size ---
    nw=0
    call ds_v3drcf2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
    & nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
    & comm2,comm3,icon)
    if(icon.ne.1000) then
        print*,'icon=',icon
        go to 9000
    endif
    allocate (z(kz1,kz2,kz3),w(nw))
    print*,'nrank,nrank2,nrank3=',nrank,nrank2,nrank3,
& ' Z-pencil x-range=',nz1b,nz1e,
& ' y-range=',nz2b,nz2e,
& ' z-range=',1,n3
    --- forward FFT ---
    idir=1
    isn=1
    call ds_v3drcf2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
    & nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
    & comm2,comm3,icon)
    if(icon.ne.0) then
        print*,'icon=',icon
        go to 9000
    endif
    --- backward FFT ---
    idir=-1
    isn=-1
    call ds_v3drcf2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
    & nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
    & comm2,comm3,icon)
```

```
        if(icon.ne.0) then
            print*,'icon=',icon
            go to 9000
        endif
c --- check result ---
    errorx=0
    do i3=1,nx3
        do i2=1,n\times2
            do i1=1,n1
                errorx=max(dabs(wc(i1,i2,i3)-
                                    x(i1,i2,i3)/n1/n2/n3),errorx)
            enddo
            enddo
        enddo
        call mpi_allreduce(errorx,errormax,1,mpi_double_precision,
        &
                            mpi_max,mpi_comm_world,ierr)
        if(nrank.eq.0)then
        print*,'num proc=',nsize
        print*,'nd2,nd3=',nsize2,nsize3
        print*,'-----(',n1,',',n2,',',n3,')-----'
        print*,'error=',errormax
    endif
C
9000 continue
    call mpi_comm_free(comm2,ierr)
    call mpi_comm_free(comm3,ierr)
    call mpi_finalize(ierr)
    stop
    end
```


## DS_V3DRCF3

Three-dimensional discrete real Fourier transforms. (mixed radices of 2, 3, 5 and 7, volumetric decomposition)

CALL DS_V3DRCF3 (X, KX1, KX2, KX1P, KX2P, KX3P, N1, N2, N3, ND1, ND2, ND3, W, NW, ISIN, ISN, COMM, ICON)

## (1) Function

The subroutine DS_V3DRCF3 performs a three-dimensional real Fourier transform or its inverse Fourier transform using a mixed radix FFT.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of $2,3,5$ and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 1 j_{j 23}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}= & \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1 r} \omega_{n 2}^{-j 2 k 2 r} \omega_{n 3}^{-j 3 k 3 r} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
& , k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
& , \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
& , \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
& , \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right) \\
& r=1 \text { or } r=-1
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j 2 j 3}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1 r} \omega_{n 2}^{j 2 k 2 r} \omega_{n 3}^{j 3 k 3 r} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right) \\
r=1 \text { or } r=-1
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality on a massively parallel machine. The global data of the three-dimensional array can be distributed among processes which are regarded as a three-dimensional grid, therefore the volumetric
decomposition allows the distribution of work more efficiently than a slabwise decomposition in an environment where massively parallel processes are available.
(2)

Parameters
X $\qquad$ Input/Output. Real data.

When the global three dimensional real data is regarded as an array $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2$, N3) virtually, the local array $X$ stores subarray of $D$ distributed along with the shape of the process grid.
If each $\mathrm{N} m$ is divisible by $\mathrm{ND} m(m=1,2,3)$ and $\mathrm{N} 1 / \mathrm{ND} 1$ is an even number, setting width parameters $\mathrm{KXmP}=\mathrm{N} m / \mathrm{NDm}$ is acceptable, then the data size settled in the subarray can be (KX1P,KX2P,KX3P) in most processes, except the edge processes of the first dimension of the process grid, which have the data size of (KX1P +2,KX2P,KX3P). Note that the array with the slightlyincreased area is necessary for every process.

If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$ for $m=2,3$, setting $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than KXmP . As for $m=1$, set an even number to KX1P such that $(\mathrm{N} 1 / 2+1) \times 2 \leq \mathrm{KX} 1 \mathrm{P} \times \mathrm{ND} 1$, when $\mathrm{KX} 1 \mathrm{P} \times \mathrm{ND} 1=\mathrm{N} 1$ is not adoptable.
The array X in each process stores the subarray of D as follows:

```
X (1:N1P, 1:N2P, 1:N3P) \(\leftarrow\) D(N1S:N1E, N2S:N2E, N3S:N3E),
N1S \(=\) KX1P \(\times\) rank1 +1 ,
\(\mathrm{N} 1 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 1, \mathrm{KX} 1 \mathrm{P} \times(\) rank1 +1\()\) ),
N1P = MAX (0, N1E - N1S + 1),
N2S \(=\) KX2P \(\times\) rank2 +1 ,
\(\mathrm{N} 2 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 2, \mathrm{KX} 2 \mathrm{P} \times(\) rank2 +1\()\) ),
\(\mathrm{N} 2 \mathrm{P}=\mathrm{MAX}(0, \mathrm{~N} 2 \mathrm{E}-\mathrm{N} 2 \mathrm{~S}+1)\),
\(\mathrm{N} 3 \mathrm{~S}=\mathrm{KX} 3 \mathrm{P} \times \operatorname{rank} 3+1\),
\(\mathrm{N} 3 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times(\) rank3 +1\())\),
N3P = MAX (0, N3E - N3S + 1),
```

where, rank1,rank2 and rank3 are coordinates of the process grid calculated as follows from the rank value, which is obtained by subroutine

```
MPI_COMM_RANK of MPI:
    rank1 = mod(rank, ND1),
    rank2 = mod(rank/ND1, ND2),
    rank3 = rank/(ND1×ND2).
```

For the real to complex transform ( $\mathrm{ISN}=1$ ), data is input; for the complex to real transform (ISN = -1), data is output.

Output/input. The real and imaginary parts of the transformed complex data.
The complex data $\mathrm{CD}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ obtained from real data $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ by Fourier transform has the complex conjugate relation so the about a half of the first dimension $(1 \sim \mathrm{~N} 1 / 2+1)$ is used to store the complex data.
(See note 2) in (3), "Comments on use.")
Regarding an array X as $\mathrm{X}(2, \mathrm{KX} 1 / 2, \mathrm{KX} 2, \mathrm{KX} 3 \mathrm{P})$, the real and imaginary parts are stored in $\mathrm{X}(1,1: \mathrm{N} 1 \mathrm{P}, 1: \mathrm{N} 2 \mathrm{P}, 1: \mathrm{N} 3 \mathrm{P})$ and $\mathrm{X}(2,1: \mathrm{N} 1 \mathrm{P}, 1: \mathrm{N} 2 \mathrm{P}, 1: \mathrm{N} 3 \mathrm{P})$ respectively, where,

```
N1S = KX1P/2 < rank1 + 1,
N1E = MIN(N1/2+1, KX1P/2 x (rank1 + 1)),
N1P = MAX(0,N1E - N1S + 1).
```

Additionally, the edge processes of the first dimension of the process grid have the increased data in $\mathrm{X}(1: 2, \mathrm{~N} 1 \mathrm{P}+1,1: \mathrm{N} 2 \mathrm{P}, 1: \mathrm{N} 3 \mathrm{P})$, when the condition KX1P $\times$ ND1 $=$ N1 is satisfied.

For the real to complex transform (ISN = 1), data is output; for the complex to real transform (ISN =-1), data is input.

This is a double precision real three-dimensional array X(KX1, KX2, KX3P).
KX1 ............ Input. The size of the first dimension of the array X. KX1 must be even.
The condition $\mathrm{KX1} \geq \mathrm{KX1P}+2$ must be satisfied if $\mathrm{KX} 1 \mathrm{P} \times \mathrm{ND} 1=\mathrm{N} 1$. Otherwise, KX1 $\geq \mathrm{KX} 1 \mathrm{P}$.

Integer(INTEGER*4).
KX2 ............ Input. The size of the second dimension of the array X ( $\geq$ KX2P). Integer(INTEGER*4).

KX1P Input. The size by which the first dimension is equally partitioned (KX1P $\times$ ND1 $\geq \mathrm{N} 1$ ). KX1P must be even. Integer(INTEGER*4).
KX2P........... Input. The size by which the second dimension is equally partitioned.(KX2P $\times$ ND $2 \geq \mathrm{N} 2$ )

Integer(INTEGER*4).
KX3P .......... Input. The size by which the third dimension is equally partitioned (KX3P $\times$ ND $3 \geq$ N3)

Integer(INTEGER*4).
N1 ............... Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N2 $\qquad$ Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
ND1 ............. Input. The number of processes by which the first dimension is partitioned.
Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
ND2 ............. Input. The number of processes by which the second dimension is partitioned.
Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
ND3 $\qquad$ Input. The number of processes by which the third dimension is partitioned.

Integer (INTEGER*4)
(See note 2) in (3), "Comments on use.")
W $\qquad$ Work area. This is double precision complex one-dimensional array W(NW).

NW $\qquad$ Input. The size of the work array $\mathrm{W}(\mathrm{NW} \geq \mathrm{MAX}(\mathrm{KX} 1 \times \mathrm{ND} 1, \mathrm{KX} 2 \mathrm{P} \times \mathrm{ND} 2$, KX3P $\times$ ND3) $\times 6$ ). It is recommended to specify a sufficiently large size for efficiency. (See note 2 ) in (2), "Comments on use.")

Integer (INTEGER*4)
ISIN $\qquad$ Input. The direction of the transformation.
$\operatorname{ISIN}=1$ for $r=1$.
ISIN $=-1$ for $r=-1$.
Integer (INTEGER*4).
ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
COMM ........ Input. The communicator indicating a set of processes on which data are distributed and by which computation is done in parallel.

Integer (INTEGER*4).
ICON $\qquad$ Output. Condition code.

See Table DS_V3DRCF3-1.
Integer (INTEGER*4).

Table DS_V3DRCF3-1 Condition codes

| Code | Meaning | Processing |
| :---: | :---: | :---: |
| 0 | No error |  |
| 25000 | Too small work area. | Processing is discontinued. |
| 30000 | N1<1, N2<1, N3<1, <br> $\mathrm{KX1}<\mathrm{KX1P}, \mathrm{KX1}$ or KX1P is not even, <br> KX2 < KX2P, N1 > KX1P $\times$ ND1, <br> N2 > KX2P $\times$ ND2, N3 > KX3P $\times$ ND3, <br> $(\mathrm{N} 1 / 2+1) \times 2>\mathrm{KX} 1 \times \mathrm{ND} 1$, <br> ISIN $\neq 1,-1$, ISN $\neq 1,-1$. |  |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |
| 30100 | ND $1 \times$ ND $2 \times$ ND3 is not equal to total processes. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1  \tag{3.1}\\
, k_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0 k 3=0}^{n 2-1} \sum_{k 1 k 2 k 3}^{n 3-1} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j_{j 2 j}}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) The results of the three-dimensional real Fourier transform has the following complex conjugate relation (indicated by ${ }^{-}$).

$$
\begin{equation*}
\alpha_{k 1 k 2 k 3}=\overline{\alpha_{n 1-k 1 n 2-k 2 n 3-k 3}} \tag{3.3}
\end{equation*}
$$

The remainder of the data is obtained from data in $k_{1}=0, \ldots, n_{1}-1, k_{2}=0, \ldots, n_{2-}$ 1 , and $k_{3}=0, \ldots, n_{3} / 2$.
3) The size of work area W

The size of the work array determines partition sizes for transferring data among nodes and calculation on each node in this routine. Setting the size NW of work area W much larger than MAX $(\mathrm{KX} 1 \times$ ND1, $\mathrm{KX} 2 \mathrm{P} \times$ ND2, $\mathrm{KX} 3 \mathrm{P} \times \mathrm{ND} 3) \times 6 \times$ (number of threads in a process) is recommended. For example, setting NW > $1,000,000$ is expected to be efficient when the assigned cache size to the process is 8 MB and the array X is partitionable by that size of the work area.
4) Parameters ND1, ND2 and ND3

When using volumetric decomposition, it is recommended to adjust ND1,ND2 and ND3 to be about the comparable value of cube root of the number of total process ND $1 \times$ ND $2 \times$ ND3 for overall efficiency of transferring data. Note that the performance of this routine may deteriorate when ND1,ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.

Additionally, this routine can be used for slabwise decomposition or 2dimensional decomposition also by setting any of ND1,ND2 or ND3 to 1, when user's program exploits specific decomposition or available shape of the process grid on a system is limited.

## 5) Consistency of parameters among processes

The parameters KX1, KX2, KX1P, KX2P, KX3P, N1, N2, N3, ND1, ND2, ND3, NW, ISIN and ISN needs to have same value respectively among all processes, otherwise the result is not guaranteed.
b. Example

Three-dimensional FFT is computed in $2 \times 2 \times 2$ processes.
** example program **
use mpi
implicit real*8 (a-h,o-z)
parameter (n1=512,n2=n1,n3=n2)
parameter (nd1=2, nd2=2, nd3=2)
parameter $(k \times 1 p=((n 1+n d 1-1) / n d 1+1) / 2 * 2)$
parameter $(k \times 1=((n 1 / 2+1)+n d 1-1) / n d 1 * 2)$
parameter $(k x 2 p=(n 2+n d 2-1) / n d 2, k \times 2=k \times 2 p)$
parameter (kx3p=(n3+nd3-1)/nd3)
parameter ( $n w=k \times 1 * k \times 2 p * k \times 3 p$ )
parameter (nwork=388)
real*8 dwork(nwork)
real* $8 \times(k \times 1, k \times 2, k \times 3 p), w(n w), w c(k \times 1, k \times 2, k \times 3 p)$
call mpi_init( ierr )
call mpi_comm_size( mpi_comm_world, nump, ierr )
call mpi_comm_rank( mpi_comm_world, nop, ierr )
nrank1=mod(nop, nd1)
nrank2=mod(nop/nd1, nd2)
nrank3=nop/(nd1*nd2)
$i x=1000$
ix=ix*nump+nop ! different seed
do i1=1, kx3p
call dvrau4(ix, x(1,1,i1),
\$ kx1*kx2,
\& dwork, nwork, icon)
enddo
do i3=1, kx3p
do i2=1, kx2p
do i1=1, kx1p
wc(i1, i2, i3) $=x(i 1, i 2, i 3)$
enddo
enddo
enddo
isin=1
isn=1
call ds_v3drcf3(x, kx1, kx2, kx1p, kx2p, kx3p,
\$ n1,n2,n3,nd1, nd2, nd3, w, nw, isin, isn,
\$ mpi_comm_world,icon)
print*,'icon=',icon
if(icon.ne.0) go to 9000

C
isn=-1
call ds_v3drcf3(x,kx1,kx2,kx1p,kx2p,kx3p,
\$ n1,n2,n3,nd1,nd2,nd3,w,nw,isin,isn,
\$ mpi_comm_world,icon)
print*,'icon=',icon
if(icon.ne.0) go to 9000
errorx=0
iof1=nrank1*kx1p
iof2=nrank2*kx2p
iof3=nrank3*kx3p
do i1=1,min(kx1p,max(n1-iof1, 0))
do i2=1, min $(k \times 2 p, \max (n 2-i o f 2,0))$
do i3=1, min(kx3p, max (n3-iof3,0))
errorx=max(dabs(wc(i1,i2,i3)-
\$ $\quad x(i 1, i 2, i 3) / n 1 / n 2 / n 3), e r r o r x)$
enddo
enddo
enddo
call mpi_allreduce(errorx,errormax,1,mpi_double_precision, \$ mpi_max,mpi_comm_world,ierr)
C
if(nop.eq.0)then
print*,'----(', n1, ', ', n2, ', ', n3, ')-----'
print*,'error=',errormax
endif
9000 continue
c
call mpi_finalize( ierr )
stop
end

## SS_V3DCFT2X

Three-dimensional discrete complex Fourier transforms. (mixed radices of 2, 3, 5 and 7, pencil decomposition, single precision)

CALL SS_V3DCFT2X (X, KX1, KX2, KX2P, KX3P, Z, KZ1, KZ2, KZ3, NZ1B, NZ1E, NZ2B, NZ2E, N1, N2, N3, W, NW, ISN, IDIR, COMM2, COMM3, ICON)
(1) Function

The subroutine SS_V3DCFT2X performs a three-dimensional complex Fourier transform or its inverse Fourier transform using a mixed radix FFT with single-precision.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of $2,3,5$ and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 11_{2 j 3}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 112 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3} & =\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
& , k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
, & k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
, & \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, & \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, & \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j 2 j}}\right\}$.

$$
\begin{gather*}
x_{j_{1 j 2} j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality using pencil decomposition. The global data of the three-dimensional array is to be distributed among a two-dimensional process grid. The local input array and output array of each process store distinct shapes by dividing the global data in different directions of pencils. This
allows the routine to omit communication for transposing back to the original distribution shape. This subroutine is the single-precision version of the DS_V3DCFT2X.
(2) Parameters

In the following, the global three dimensional complex data is regarded as an array $\mathrm{D}(\mathrm{N} 1$, N2, N3) virtually, and the process grid is regarded as ND2 $\times$ ND3 shape.

X . $\qquad$ Input when IDIR $=1$. Complex data

The local array X stores a subarray of D, which corresponds to a columnwise decomposed part of the global array by dividing the second and third dimensions by ND2 and ND3 respectively.

If each $\mathrm{N} m$ is divisible by $\mathrm{ND} m$ ( $m=2,3$ ), setting a width parameter $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m$ is recommended, then the sizes of the subarray can be (N1,KX2P,KX3P) uniformly.
If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$, setting $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than KXmP .

The array X in each process stores the subarray of D as follows:
X(1:N1, 1:NX2P, 1:NX3P) $\leftarrow \mathrm{D}(1: N 1$, NX2B:NX2E, NX3B:NX3E)
NX2B $=$ KX2P $\times$ rank2 +1
$\mathrm{NX} 2 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 2, \mathrm{KX} 2 \mathrm{P} \times($ rank $2+1))$
NX2P $=$ MAX $(0$, NX2E - NX2B +1$)$
NX3B $=$ KX3P $\times$ rank3 +1
$\mathrm{NX} 3 \mathrm{E}=\operatorname{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times(\operatorname{rank} 3+1))$
NX3P $=\operatorname{MAX}(0$, NX3E $-N X 3 B+1)$,
where the rank2 and rank3 are ranks of the process in the communicator COMM2 and COMM3 respectively, which are obtained by subroutine MPI_COMM_RANK of MPI.

The input values are not retained after the calculation.
......... Output when IDIR = - 1. Transformed complex data.
Resultant transformed three dimensional data of $\mathrm{D}(\mathrm{N} 1, \mathrm{~N} 2, \mathrm{~N} 3)$ are stored into the array X in the same distributed way as stated above.

This is a single precision complex three-dimensional array X(KX1, KX2, NX3P).
KX1 ............ Input. The size of the first dimension of the array $\mathrm{X}(\geq \mathrm{N} 1)$. Integer(INTEGER*4).
KX2 ............ Input. The size of the second dimension of the array $X(\geq$ NX2P $)$. Integer(INTEGER*4).
KX2P
Input. The size by which the second dimension of D is equally partitioned when the data are stored in the array X . ( $\mathrm{KX} 2 \mathrm{P} \times \mathrm{ND} 2 \geq \mathrm{N} 2$ ) Integer(INTEGER*4).
KX3P .......... Input. The size by which the third dimension of D is equally partitioned when the data are stored in the array X . (KX3P $\times \mathrm{ND} 3 \geq \mathrm{N} 3$ )

Integer(INTEGER*4).
Z $\qquad$ Output when IDIR = 1. Transformed complex data.

The local array Z stores a subarray of D , which corresponds to a columnwise decomposed part of the global array by dividing the first and second dimensions by ND2 and ND3 respectively.
The array Z in each process stores the subarray of D as follows:

$$
\begin{aligned}
& \text { Z(1:NZ1P, 1:NZ2P, 1:N3) ↔ D(NZ1B:NZ1E, NZ2B:NZ2E, 1:N3) } \\
& \text { NZ1P = MAX }(0, \text { NZ1E - NZ1B + 1) } \\
& \text { NZ2P = MAX }(0, \text { NZ2E }- \text { NZ2B + 1). }
\end{aligned}
$$

..........Input when IDIR $=-1$. Complex data.
The local array Z stores the subarray of D in the same distributed way as stated above.

This is a single precision complex three-dimensional array Z(KZ1, KZ2, KZ3).
KZ1 ............. Output when $\mathrm{NW}=0$. The recommended size for the first dimension of the array Z.

Input when $N W \neq 0$. The size of the first dimension of the array $Z(\geq$ NZ1ENZ1B+1).

Integer(INTEGER*4).
KZ2 ............. Output when NW $=0$. The recommended size for the second dimension of the array Z.

Input when $N W \neq 0$. The size of the second dimension of the array $Z(\geq$ NZ2E-NZ2B+1). Integer(INTEGER*4).
KZ3 ............. Output when NW $=0$. The recommended size of the third dimension of the array Z.

Input when $N W \neq 0$. The size of the third dimension of the array $Z(\geq N 3)$. Integer(INTEGER*4).

NZ1B $\qquad$ Output. The starting index for the first dimension of the global array D. Integer(INTEGER*4).
NZ1E $\qquad$ Output. The ending index for the first dimension of the global array D. NZ1B and NZ1E indicate which portion of the first dimension within the global array is stored in the local array Z . Integer(INTEGER*4).

NZ2B $\qquad$ Output. The starting index for the second dimension of the global array D. Integer(INTEGER*4).

NZ2E $\qquad$ Output. The ending index for the second dimension of the global array D. NZ2B and NZ2E indicate which portion of the second dimension within the global array is stored in the local array Z . Integer(INTEGER*4).
N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
$\qquad$ Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
W $\qquad$ Work area. This is single precision complex one-dimensional array W(NW).
NW $\qquad$ Input / Output. The size of the work array W.
When NW $=0$ specified, the recommended sizes of NW, KZ1, KZ2, and KZ3 are set respectively, and index information is set to NZ1B, NZ1E, NZ2B, and NZ2E.
Integer (INTEGER*8). (See note 4) in (3), "Comments on use.")
ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN $=1$ for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
IDIR $\qquad$ Input. The direction of transform between arrays is indicated.
IDIR $=1$ for the transform from the array X to the array Z .
IDIR $=-1$ for the transform from the array Z to the array X .
Integer (INTEGER*4).
COMM2 ..... Input. The MPI communicator that represents a set of processes whose size of the process group is ND2, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 2) in (3), "Comments on use.")

Integer (INTEGER*4).
COMM3 ..... Input. The MPI communicator that represents a set of processes whose size of the process group is ND3, which is obtained by MPI_COMM_SIZE, in the process shape ND $2 \times$ ND3. (See note 2) in (3), "Comments on use.") Integer (INTEGER*4).
ICON $\qquad$ Output. Condition code.
See Table DS_V3DCFT2X-1. Integer (INTEGER*4).

Table SS_V3DCFT2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |

Table SS_V3DCFT2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 1000 | NW $=0$ is specified. | The recommended sizes of <br> NW, KZ1, KZ2, and KZ3 <br> are set. <br> Index information is set to <br> NZ1B, NZ1E, NZ2B, and <br> NZ2E. <br> There is no output to array X <br> or Z. |
| 25000 | Too small work area. | Processing is discontinued. |
| 30000 | N1<1, N2<1, N3<1, KX1 < N1, KX2 < NX2P, <br> N2 $>$ KX2P $\times$ ND2, N3 $>$ KX3P $\times$ ND3, or the <br> value of ISN or IDIR is incorrect. |  |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 33=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1 \\
, k_{3}=0,1, \ldots, n_{3}-1 \tag{3.1}
\end{gather*}
$$

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1  \tag{3.2}\\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j_{j 2 j}}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) Process shape ND2 and ND3

Note that the performance of this routine may deteriorate when ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.
3) Consistency of parameters among processes

The parameters KX2P,KX3P,N1,N2,N3,NW,ISN, and IDIR needs to have same value respectively among all processes, otherwise the result is not guaranteed.
4) The size of work area W

The size of the work array NW needs to be about twice the size of array X or Z to be used as send/receive buffers for MPI communication inside the routine. Note that the parameter NW is an 8-byte integer type.
b. Example

Three-dimensional FFT is computed in $2 \times 3$ processes.
c --- prepare test-data --
$n \times 2=\min (k x 2 p, \max (n 2-n r a n k 2 * k x 2 p, 0))$
$n \times 3=\min (k \times 3 p, \max (n 3-n r a n k 3 * k x 3 p, 0))$
ix=1000
ix=ix+nrank ! different seed
do i3=1, nx3

```
        do i2=1,n\times2
            call ranu2(ix,x(1,i2,i3),2*n1,icon)
            do i1=1,n1
                wc(i1,i2,i3)=x(i1,i2,i3)
            enddo
        enddo
        enddo
    --- inquire necessary size ---
    nw=0
    call ss_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
& nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
& comm2,comm3,icon)
    if(icon.ne.1000) then
        print*,'icon=',icon
        go to 9000
    endif
    allocate (z(kz1,kz2,kz3),w(nw))
    print*,'nrank,nrank2,nrank3=',nrank,nrank2,nrank3,
& ' Z-pencil x-range=',nz1b,nz1e,
& ' y-range=',nz2b,nz2e,
& ' z-range=',1,n3
    --- forward FFT ---
    idir=1
    isn=1
    call ss_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
& nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
& comm2,comm3,icon)
    if(icon.ne.0) then
        print*,'icon=',icon
        go to 9000
    endif
    --- backward FFT ---
    idir=-1
    isn=-1
    call ss_v3dcft2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
& nz1b,nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
& comm2,comm3,icon)
    if(icon.ne.0) then
        print*,'icon=',icon
        go to 9000
    endif
    --- check result ---
    errorx=0
    do i3=1,nx3
    do i2=1,n\times2
        do i1=1,n1
            errorx=max(cabs(wc(i1,i2,i3)-
&
                                    x(i1,i2,i3)/n1/n2/n3),errorx)
                enddo
        enddo
    enddo
    call mpi_allreduce(errorx,errormax,1,mpi_real,
& mpi_max,mpi_comm_world,ierr)
```

```
        if(nrank.eq.0)then
            print*,'num proc=',nsize
            print*,'nd2,nd3=',nsize2,nsize3
            print*,'-----(',n1,',',n2,',',n3,')-----'
            print*,'error=',errormax
            endif
c
    9 0 0 0 ~ c o n t i n u e
        call mpi_comm_free(comm2,ierr)
        call mpi_comm_free(comm3,ierr)
        call mpi_finalize(ierr)
        stop
        end
```


## SS_V3DRCF2X

Three-dimensional discrete real Fourier transforms. (mixed radices of 2, 3, 5 and 7, pencil decomposition, single precision)

CALL SS_V3DRCF2X (X, KX1, KX2, KX2P, KX3P, Z, KZ1, KZ2, KZ3, NZ1B, NZ1E, NZ2B, NZ2E, N1, N2, N3, W, NW, ISN, IDIR, COMM2, COMM3, ICON)
(1) Function

The subroutine SS_V3DRCF2X performs a three-dimensional real Fourier transform or its inverse Fourier transform using a mixed radix FFT with single-precision.

The size of each dimension of three-dimensional arrays $\left(n_{1}, n_{2}, n_{3}\right)$ can be a product of the powers of $2,3,5$ and 7 .
a. The three-dimensional Fourier transform

When $\left\{x_{j 11_{2 j 3}}\right\}$ is input, the transform defined by (1.1) below is calculated to obtain $\left\{n_{1} n_{2} n_{3} \alpha_{k 112 k 3}\right\}$.

$$
\begin{align*}
n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3} & =\sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 3=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, & k_{1}=0,1, \ldots, n_{1}-1 \\
& , k_{2}=0,1, \ldots, n_{2}-1 \\
& , k_{3}=0,1, \ldots, n_{3}-1  \tag{1.1}\\
& , \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
& , \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, & \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{align*}
$$

b. The three-dimensional Fourier inverse transform

When $\left\{\alpha_{k 1 k 2 k 3}\right\}$ is input, the transform defined by (1.2) below is calculated to obtain $\left\{x_{j 1 j_{j 2 j}}\right\}$.

$$
\begin{gather*}
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3} \\
, j_{1}=0,1, \ldots, n_{1}-1 \\
, j_{2}=0,1, \ldots, n_{2}-1 \\
, j_{3}=0,1, \ldots, n_{3}-1  \tag{1.2}\\
, \omega_{n 1}=\exp \left(2 \pi i / n_{1}\right) \\
, \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right) \\
, \omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
\end{gather*}
$$

This subroutine provides an efficient and scalable 3D FFT functionality using pencil decomposition. The global data of the three-dimensional array is to be distributed among a two-dimensional process grid. The local input array and output array of each process store distinct shapes by dividing the global data in different directions of pencils. This
allows the routine to omit communication for transposing back to the original distribution shape. This subroutine is the single-precision version of the DS_V3DRCF2X.
(2) Parameters

In the following, the global three dimensional real data is regarded as an array $\operatorname{DR}(\mathrm{N} 1, \mathrm{~N} 2$, N3), and the transformed global three dimensional complex data is regarded as an array $\mathrm{DC}(\mathrm{N} 1 / 2+1, \mathrm{~N} 2, \mathrm{~N} 3)$ virtually, and the process grid is regarded as $\mathrm{ND} 2 \times \mathrm{ND} 3$ shape.

X . $\qquad$ Input when IDIR = 1. Real data.
The local array X stores a subarray of DR , which corresponds to a columnwise decomposed part of the global array by dividing the second and third dimensions by ND2 and ND3 respectively.
If each $\mathrm{N} m$ is divisible by $\mathrm{NDm}(m=2,3)$, setting width parameters $\mathrm{KXmP}=\mathrm{N} m / \mathrm{ND} m$ is recommended, then the data size settled in the subarray can be (N1,KX2P,KX3P) uniformly.
If $\mathrm{N} m$ is not divisible by $\mathrm{ND} m$, setting $\mathrm{KX} m \mathrm{P}=\mathrm{N} m / \mathrm{ND} m+1$ is recommended, then the sizes of the subarrays that are assigned to the edge of the process grid are less than $\mathrm{KX} m \mathrm{P}$.

The array X in each process stores the subarray of DR as follows:
X(1:N1, 1:NX2P, 1:NX3P) $\leftarrow$ DR(1:N1, NX2B:NX2E, NX3B:NX3E),
$\mathrm{NX} 2 \mathrm{~B}=\mathrm{KX} 2 \mathrm{P} \times$ rank2 +1
$\mathrm{NX} 2 \mathrm{E}=\operatorname{MIN}(\mathrm{N} 2, \mathrm{KX} 2 \mathrm{P} \times(\operatorname{rank} 2+1))$
NX2P $=\operatorname{MAX}(0$, NX2E - NX2B +1$)$
NX3B $=$ KX3P $\times$ rank3 +1
$\mathrm{NX} 3 \mathrm{E}=\mathrm{MIN}(\mathrm{N} 3, \mathrm{KX} 3 \mathrm{P} \times($ rank3 +1$))$
NX3P $=\operatorname{MAX}(0$, NX3E - NX3B +1$)$,
where, rank2 and rank3 are the ranks of the process in the communicator COMM2 and COMM3 respectively, which are obtained by subroutine MPI_COMM_RANK of MPI.

The input values are not retained after the calculation.
......... Output when IDIR $=-1$. Transformed real data.
Resultant transformed three dimensional real data from DC(N1/2+1, N2, N3) are stored into the array X in the same distributed way as stated above.

This is a single precision real three-dimensional array X(KX1, KX2, KX3P).
KX1 ............ Input. The size of the first dimension of the array X. KX1 must be even. ( $\geq \mathrm{N} 1$ ). Integer(INTEGER*4).

KX2 ............ Input. The size of the second dimension of the array $X(\geq N X 2 P)$.
Integer(INTEGER*4).
KX2P. $\qquad$ Input. The size by which the second dimension of DR is equally partitioned. (KX2P $\times \mathrm{ND} 2 \geq \mathrm{N} 2$ ) Integer(INTEGER*4).
KX3P .......... Input. The size by which the third dimension of DR is equally partitioned. (KX3P $\times$ ND $3 \geq$ N3)

Integer(INTEGER*4).
Z $\qquad$ Output when IDIR = 1. Transformed complex data.

The complex data obtained from real data DR(N1, N2, N3) by Fourier transform has the complex conjugate relation so the about a half of the first dimension is used to store. The local array Z stores a subarray of $\mathrm{DC}(\mathrm{N} 1 / 2+1$, N2, N3), which corresponds to a columnwise decomposed part of the global array by dividing the first and second dimensions by ND2 and ND3 respectively.

The array Z in each process stores the subarray of DC as follows: Z(1:NZ1P, 1:NZ2P, 1:N3) $\leftarrow$ DC(NZ1B:NZ1E, NZ2B:NZ2E, 1:N3)
NZ1P = MAX(0, NZ1E - NZ1B + 1)
NZ2P $=$ MAX $(0$, NZ2E - NZ2B +1$)$.
..........Input when IDIR $=-1$. Complex data.
The local array Z stores the subarray of DC in the same distributed way as stated above.

This is a single precision complex three-dimensional array Z(KZ1, KZ2, KZ3).
(See note 2) in (3), "Comments on use.")
KZ1 ............. Output when $\mathrm{NW}=0$. The recommended size for the first dimension of the array Z.

Input when $N W \neq 0$. The size of the first dimension of the array $Z(\geq$ NZ1E NZ1B+1).

Integer(INTEGER*4)
KZ2 ............. Output when NW $=0$. The recommended size for the second dimension of the array Z.
Input when $N W \neq 0$. The size of the second dimension of the array $Z(\geq$ NZ2E-NZ2B+1).

Integer(INTEGER*4).
KZ3 $\qquad$ Output when NW $=0$. The recommended size of the third dimension of the array Z.

Input when $\mathrm{NW} \neq 0$. The size of the third dimension of the array $\mathrm{Z}(\geq \mathrm{N} 3)$. Integer(INTEGER*4).

NZ1B $\qquad$ Output. The starting index for the first dimension of the global array DC. Integer(INTEGER*4).

NZ1E .......... Output. The ending index for the first dimension of the global array DC. NZ1B and NZ1E indicate which portion of the first dimension within the global array is stored in the local array Z .

Integer(INTEGER*4).
NZ2B .......... Output. The starting index for the second dimension of the global array DC. Integer(INTEGER*4).

NZ2E $\qquad$ Output. The ending index for the second dimension of the global array DC. NZ2B and NZ2E indicate which portion of the second dimension within the global array is stored in the local array Z .

Integer(INTEGER*4)
N1 $\qquad$ Input. The length $n_{1}$ of data in the first dimension of the three- dimensional array to be transformed.
$n_{1}$ must be a value that can be a product of the powers of $2,3,5$ and 7 . Integer (INTEGER*4)
N2 ............... Input. The length $n_{2}$ of data in the second dimension of the three- dimensional array to be transformed.
$n_{2}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
N3 $\qquad$ Input. The length $n_{3}$ of data in the third dimension of the three- dimensional array to be transformed.
$n_{3}$ must be a value that can be a product of the powers of $2,3,5$ and 7 .
Integer (INTEGER*4)
W $\qquad$ Work area. This is a single precision complex one-dimensional array W(NW).

NW ............. Input / Output. The size of the work array W.
When NW $=0$ specified, the recommended sizes of NW, KZ1, KZ2, and KZ3 are set respectively, and the index information is set to NZ1B, NZ1E, NZ2B, and NZ2E.
Integer (INTEGER*8). (See note 5) in (3), "Comments on use.")
ISN $\qquad$ Input. Either the transform or the inverse transform is indicated.

ISN = 1 for the transform.
ISN $=-1$ for the inverse transform.
Integer (INTEGER*4).
IDIR $\qquad$ Input. The direction of transform between arrays is indicated.
IDIR = 1: from the real array X to the complex array Z .
IDIR $=-1$ : from the complex array Z to the real array X .
Integer (INTEGER*4).
COMM2 ..... Input. The MPI communicator that represents a set of processes whose size of the process group is ND2, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 3) in (3), "Comments on use.")
Integer (INTEGER*4).
COMM3 ..... Input. The MPI communicator that represents a set of processes whose size of the process group is ND3, which is obtained by MPI_COMM_SIZE, in the process shape ND2 $\times$ ND3. (See note 3) in (3), "Comments on use.") Integer (INTEGER*4).
ICON $\qquad$ Output. Condition code.
See Table DS_V3DRCF2X-1. Integer (INTEGER*4).

Table SS_V3DRCF2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |

Table SS_V3DRCF2X-1 Condition codes

| Code | Meaning | Processing |
| :--- | :--- | :--- |
| 0 | No error |  |
| 1000 | NW $=0$ is specified | The recommended sizes of <br> NW, KZ1, KZ2 and KZ3 are <br> set. <br> Index information is set to <br> NZ1B, NZ1E, NZ2B, and <br> NZ2E. <br> There is no output to array X <br> or Z. |
| 25000 | Too small work area. |  |
| 30000 | N1<1, N2<1, N3<1, <br> KX1 $<$ N1, KX1 is not even, <br> KX2 $<$ NX2P, <br> N2 $>$ KX2P $\times$ ND2, N3 $>$ KX3P $\times$ ND3, <br> $(N 1 / 2+1) \times 2>$ KZ1 $\times$ ND2, <br> ISN $\neq 1,-1$, IDIR $\neq 1, ~-1$. | Processing is discontinued. |
| 30008 | The order of the transform is not radix 2/3/5/7. |  |

(3) Comments on use
a. Notes

1) General definition of a Fourier transform

The three-dimensional discrete complex Fourier transform and its inverse transform can generally be defined as in (3.1) and (3.2).

$$
\begin{gather*}
\alpha_{k 1 k 2 k 3}=\frac{1}{n_{1} n_{2} n_{3}} \sum_{j 1=0}^{n 1-1} \sum_{j 2=0}^{n 2-1} \sum_{j 33=0}^{n 3-1} x_{j 1 j 2 j 3} \omega_{n 1}^{-j 1 k 1} \omega_{n 2}^{-j 2 k 2} \omega_{n 3}^{-j 3 k 3} \\
, k_{1}=0,1, \ldots, n_{1}-1 \\
, k_{2}=0,1, \ldots, n_{2}-1  \tag{3.1}\\
, k_{3}=0,1, \ldots, n_{3}-1
\end{gather*}
$$

$$
x_{j 1 j 2 j 3}=\sum_{k 1=0}^{n 1-1} \sum_{k 2=0}^{n 2-1} \sum_{k 3=0}^{n 3-1} \alpha_{k 1 k 2 k 3} \omega_{n 1}^{j 1 k 1} \omega_{n 2}^{j 2 k 2} \omega_{n 3}^{j 3 k 3}
$$

$$
\begin{equation*}
, j_{1}=0,1, \ldots, n_{1}-1 \tag{3.2}
\end{equation*}
$$

$$
, j_{2}=0,1, \ldots, n_{2}-1
$$

$$
, j_{3}=0,1, \ldots, n_{3}-1
$$

where, $\omega_{n 1}=\exp \left(2 \pi i / n_{1}\right), \omega_{n 2}=\exp \left(2 \pi i / n_{2}\right)$,

$$
\omega_{n 3}=\exp \left(2 \pi i / n_{3}\right)
$$

This subroutine calculates $\left\{n_{1} n_{2} n_{3} \alpha_{k 1 k 2 k 3}\right\}$ or $\left\{x_{j 1 j_{j 2 j}}\right\}$ corresponding to the left-hand-side term of (3.1) or (3.2), respectively. Normalization of the results may be required.
2) The results of the three-dimensional real Fourier transform has the following complex conjugate relation (indicated by ${ }^{-}$).

$$
\begin{equation*}
\alpha_{k 1 k 2 k 3}=\overline{\alpha_{n 1-k 1 n 2-k 2 n 3-k 3}} \tag{3.3}
\end{equation*}
$$

The remainder of the data is obtained from data in $k_{1}=0, \ldots, n_{1}-1, k_{2}=0, \ldots, n_{2-}$ 1 , and $k_{3}=0, \ldots, n_{3} / 2$.
3) Process shape ND2 and ND3

Note that the performance of this routine may deteriorate when ND1,ND2 and ND3 do not match the shape of the executing process grid on a system which can specify the shape of the process grid. Refer to the Job Operation Software manual whether the system can assign a shape of the process grid.
4) Consistency of parameters among processes

The parameters KX2P, KX3P, N1, N2, N3, NW, ISN and IDIR needs to have same value respectively among all processes, otherwise the result is not guaranteed.
5) The size of work area W

The size of the work array NW needs to be about twice the size of array X or Z to be used as send/receive buffers for MPI communication inside the routine. Note that the parameter NW is an 8-byte integer type.
b. Example

Three-dimensional FFT is computed in $2 \times 3$ processes.

```
    ** example program **
    use mpi
    implicit real (a-h,o-z)
    parameter (n1=512,n2=n1,n3=n2)
    parameter (nd2=2,nd3=3)
    parameter (n1c=n1/2+1,kx1=n1c*2)
    parameter (kx2p=((n2+nd2-1)/nd2),kx2=kx2p)
    parameter (kx3p=((n3+nd3-1)/nd3))
    integer comm2,comm3
    integer*8 nw
    real x(kx1,kx2,kx3p),wc(kx1,kx2,kx3p)
    complex,allocatable :: z(:,:,:)
    real,allocatable :: w(:)
    --- prepare sub-communicator ---
    call mpi_init(ierr)
    call mpi_comm_size(mpi_comm_world, nsize, ierr)
    call mpi_comm_rank(mpi_comm_world, nrank, ierr)
    ncolory=nrank/nd2
    call mpi_comm_split(mpi_comm_world,ncolory,nrank,
    & comm2,ierr)
    call mpi_comm_size(comm2, nsize2, ierr)
    call mpi_comm_rank(comm2, nrank2, ierr)
    ncolorz=mod(nrank,nd2)
    call mpi_comm_split(mpi_comm_world,ncolorz,nrank,
```

                    comm3,ierr)
        call mpi_comm_size(comm3, nsize3, ierr)
        call mpi_comm_rank(comm3, nrank3, ierr)
        if(nsize.ne.nd2*nd3 .or. nsize2.ne.nd2 .or.
    \& nsize3.ne.nd3) then
        print*,'nsize=',nsize,nsize2,nsize3
        go to 9000
        endif
    --- prepare test-data ---
    \(n \times 2=\min (k \times 2 p, \max (n 2-n r a n k 2 * k x 2 p, 0))\)
    \(n \times 3=m i n(k x 3 p, \max (n 3-n r a n k 3 * k x 3 p, 0))\)
    ix=1000
    ix=ix+nrank ! different seed
    do i3=1,nx3
        do i2=1, nx 2
            call ranu2(ix, x(1,i2,i3),n1,icon)
            do i1=1,n1
                \(w c(i 1, i 2, i 3)=x(i 1, i 2, i 3)\)
            enddo
        enddo
    enddo
    --- inquire necessary size ---
    nw=0
    call ss_v3drcf \(2 x(x, k x 1, k x 2, k x 2 p, k x 3 p, z, k z 1, k z 2, k z 3\),
    \& nz1b, nz1e, nz2b, nz2e, n1, n2, n3, w, nw, isn, idir,
\& comm2, comm3,icon)
if(icon.ne.1000) then
print*,'icon=',icon
go to 9000
endif
allocate (z(kz1,kz2,kz3),w(nw))
print*,'nrank, nrank2, nrank3=', nrank, nrank2, nrank3,
\& ' Z-pencil x-range=',nz1b,nz1e,
\& ' y-range=',nz2b,nz2e,
\& ' z-range=',1,n3
--- forward FFT ---
idir=1
isn=1
call ss_v3drcf2x(x,kx1,kx2,kx2p,kx3p,z,kz1,kz2,kz3,
\& nz1b, nz1e,nz2b, nz2e, n1, n2, n3,w, nw, isn,idir,
\& comm2, comm3,icon)
if(icon.ne.0) then
print*,'icon=',icon
go to 9000
endif
--- backward FFT ---
idir=-1
isn=-1
call ss_v3drcf $2 x(x, k x 1, k x 2, k x 2 p, k x 3 p, z, k z 1, k z 2, k z 3$,
\& nz1b, nz1e,nz2b,nz2e,n1,n2,n3,w,nw,isn,idir,
\& comm2, comm3,icon)
if(icon.ne.0) then
print*,'icon=',icon

```
            go to 9000
        endif
c
    &
        --- check result ---
        errorx=0
        do i3=1,nx3
            do i2=1,n\times2
            do i1=1,n1
                errorx=max(abs(wc(i1,i2,i3)-
        enddo
            enddo
        enddo
        call mpi_allreduce(errorx,errormax,1,mpi_real,
        & mpi_max,mpi_comm_world,ierr)
        if(nrank.eq.0)then
            print*,'num proc=',nsize
            print*,'nd2,nd3=',nsize2,nsize3
            print*,'-----(',n1,',',n2,',',n3,')-----'
            print*,'error=',errormax
        endif
c
9000 continue
    call mpi_comm_free(comm2,ierr)
    call mpi_comm_free(comm3,ierr)
    call mpi_finalize(ierr)
    stop
    end
```

Appendix

## Appendix A

[1] Markus Hegland
Block Algorithms for FFTs on Vector and Parallel Computers. PARCO 93, Grenoble, 1993.
[2] Charles Van Loan
Computational Frameworks for the Fast Fourier Transform, SIAM, 1992.

## Index

M
mixed radices
$8,13,21,27,33,41,48,56$

## T

Three-dimensional discrete complex Fourier transforms.....................................8, 13, 21, 48
Three-dimensional discrete real Fourier transforms $27,33,41,56$

